Constraint Programming II: Solving CPs using Propagation and Basic Search

Slides draw upon material from: 6.034 notes, by Tomas Lozano Perez; AIMA, by Stuart Russell & Peter Norvig; Constraint Processing, by Rina Dechter. Brian C. Williams
Enrique Fernandez
16.410/413
October 28th, 2015

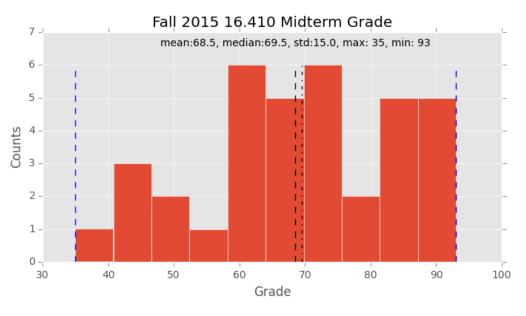
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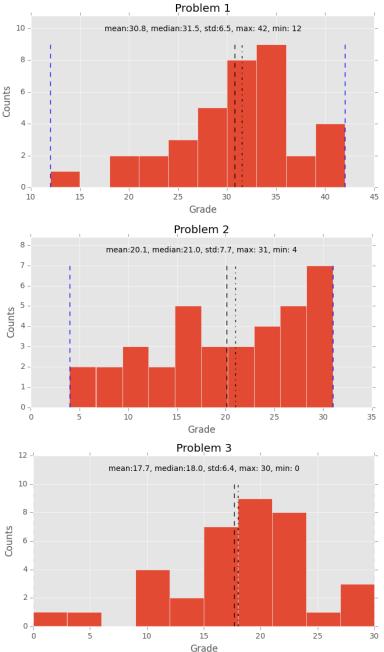
Assignments

- Remember:
 - Problem Set #6: Out today. Due next Wednesday
 - Project Part 1 (16.413): Due on Nov 6th
- Reading:
 - Today and Monday:
 [AIMA] Ch. 6.2-5; Constraint Satisfaction.
- To Learn More: Constraint Processing, by Rina Dechter.
 - Ch. 5: General Search Strategies: Look-Ahead.
 - Ch. 6: General Search Strategies: Look-Back.
 - Ch. 7: Stochastic Greedy Local Search.

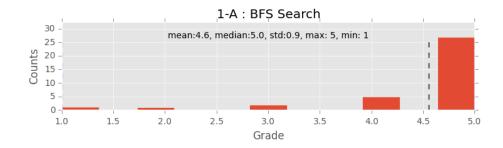
Midterm results

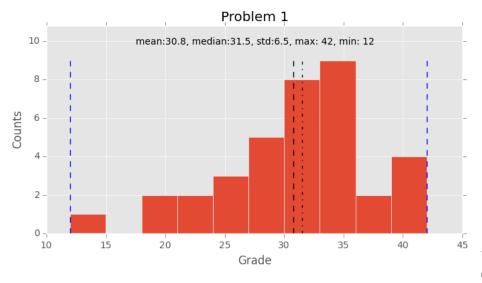
	P1	P2	P	3	Total
Max		42	31	30	93
Min		12	4	0	35
Avg		31	20	18	69
Median		32	21	18	70
Std		6.45	7.73	6.36	14.98
# 0		0	0	1	0

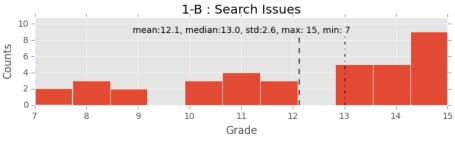


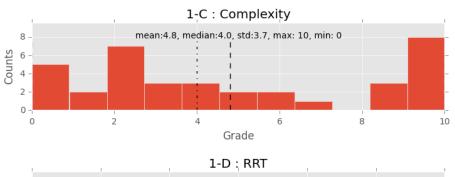


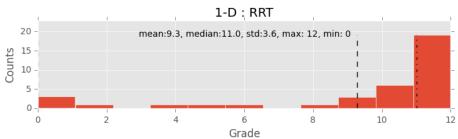
	1-A	1-B	1-C		1-D	P1 (Total)
Max	5	1	5	10	12	42
Min	1		7	0	0	12
Avg	5	1	2	5	9	31
Median	5	1	.3	4	11	. 32
Std	0.94	2.6	1 3	3.73	3.56	6.45
# 0s	0		0	5	1	. 0



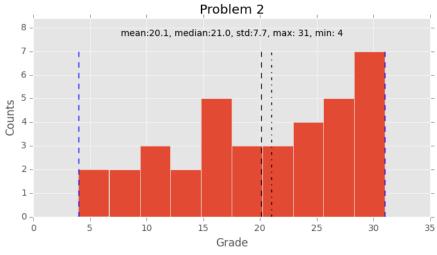


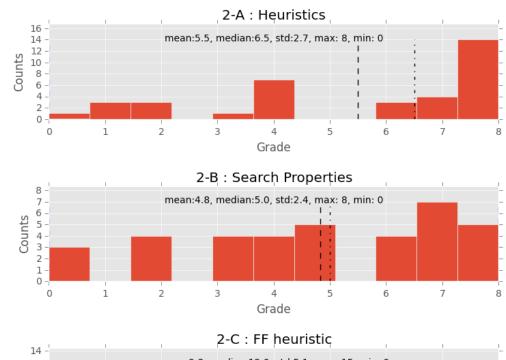


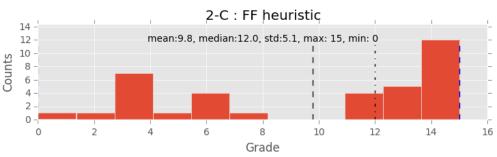




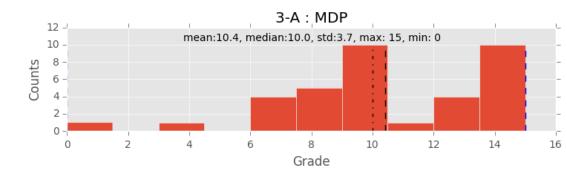
	2-A	2-B	2	2-C	P2 (Total)
Max		8	8	15	31
Min		0	0	0	4
Avg		6	5	10	20
Median		7	5	12	21
Std		2.66	2.43	5.14	7.73
# 0s		1	3	1	0

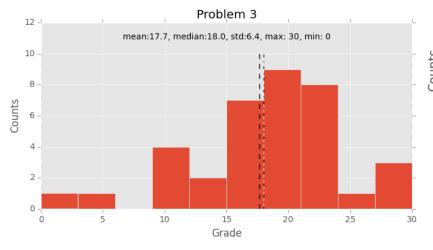


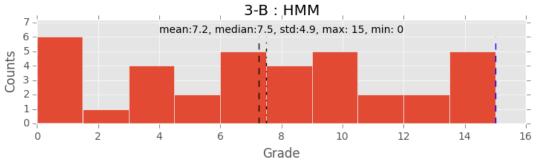




	3-A	3-B	P3	(Total)
Max		15	15	30
Min		0	0	0
Avg		10	7	18
Median		10	8	18
Std		3.68	4.92	6.36
# 0s		1	6	1







Constraint Problems are Everywhere

7	5		9		3			6
			4	5				3
6	2			9		8		
	1	5				2	3	
		9		1			7	5
3				8	4			
9			6		1		5	7

(a) Sudoku Puzzle

7	5	8	9	2	3	1	4	6
2	4	3	1	6	7	5	9	8
1	9	6	4	5	8	7	2	3
6	2	7	3	9	5	8	1	4
8	1	5	7	4	6	2	3	9
4	3	9	8	1	2	6	7	5
3	7	1	5	8	4	9	6	2
5	6	4	2	7	9	3	8	1
9	8	2	6	3	1	4	5	7

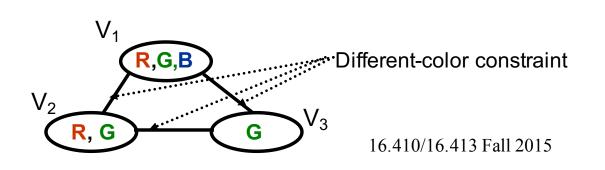
(b) The Solution

Constraint Satisfaction Problems (CSP)

Input: A Constraint Satisfaction Problem is a triple <V,D,C>, where:

- V is a set of variables V_i ,
- D is a set of variable domains,
 - The domain of variable V_i is denoted D_i,
- C = is a set of constraints on assignments to V,
 - Each constraint C_i = <S_i,R_i> specifies allowed variable assignments,
 - S_i the constraint's scope, is a subset of variables V,
 - R_i the constraint's relation, is a set of assignments to S_i.

Output: A full assignment to V, from elements of V's domain, such that all constraints in C are satisfied.



V?
$$V = \{V1, V2, V3\}$$

 D_1 ? $D_1 = \{R,G,B\}$
 C_{12} ? $C_{12} = \{, , , \}$

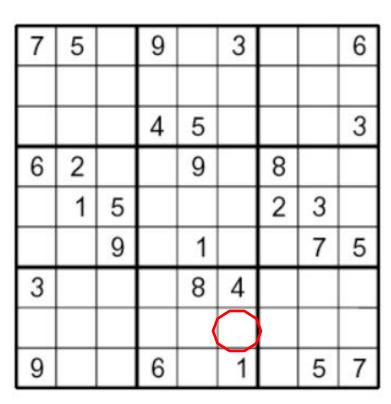
Constraint Modeling (Programming) Languages

Features Declarative specification of the problem that separates the formulation and the search strategy.

Example: Constraint Model of the Sudoku Puzzle in

Number Jack (http://4c110.ucc.ie/numberjack/home).

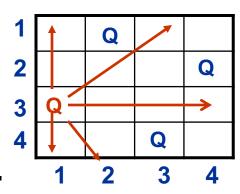
Constraint Problems are Everywhere



(a) Sudoku Puzzle

N-Queens

Place queens so that no queen can attack another.



Encoding

- Assume one queen per column.
- Determine what row each queen should be in.

Variables Q_1, Q_2, Q_3, Q_4 .

Domains {1, 2, 3, 4}.

Constraints $Q_i \iff Q_j$ "On different rows".

 $|Q_i - Q_j| \iff |i-j|$ "Stay off the diagonals".

Example $C_{1,2} = \{(1,3) \ (1,4) \ (2,4) \ (3,1) \ (4,1) \ (4,2)\}.$

Outline

- Arc-consistency and constraint propagation.
- Analysis of constraint propagation.
- Solving CSPs using search.

Good News / Bad News

Good News

- Very general & interesting family of problems.
- Problem formulation used extensively in autonomy and decision making applications.

Bad News

Includes NP-Hard (intractable ?) problems.

Algorithmic Design Paradigm

Solving CSPs involves a combination of:

1. Inference

- Solves partially by eliminating values that can't be part of any solution (constraint propagation).
- Makes implicit constraints explicit.

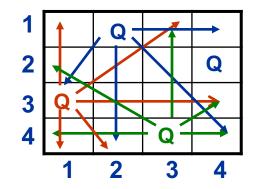
2. Search

Tries alternative assignments against constraints.

N-Queens

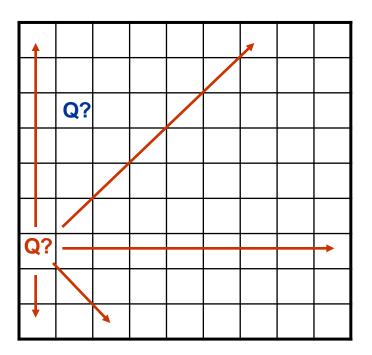
Inference

Eliminate values that can't be part of any solution



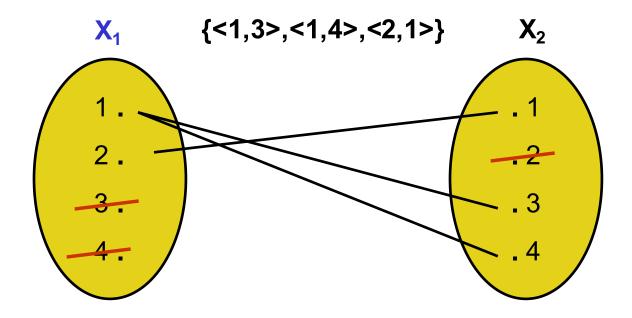
Search

Try alternative assignments against constraints



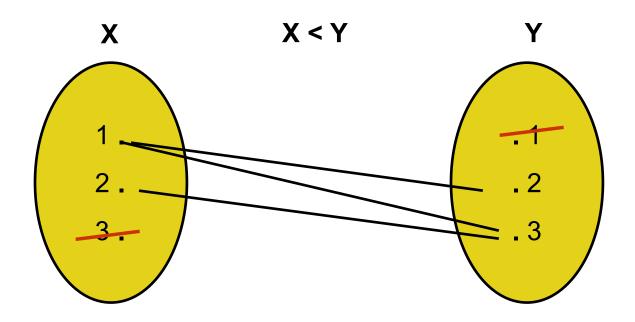
Arc Consistency

Idea: Eliminate values of a variable domain that can <u>never satisfy</u> a specified constraint (an arc).

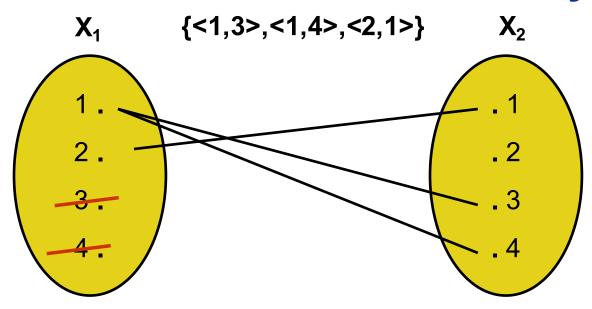


Definition: arc $\langle x_i, x_j \rangle$ is arc consistent if $\langle x_i, x_j \rangle$ and $\langle x_j, x_i \rangle$ are directed arc consistent.

Arc Consistency



Directed Arc Consistency



Definition: arc $\langle x_i, x_i \rangle$ is directed arc consistent if

- for every a_i in D_i,
 - there exists some a_i in D_i such that
 - assignment <a_i,a_j> satisfies constraint C_{ij}
- \forall $a_i \in D_i$, \exists $a_j \in D_j$ such that $\langle a_i, a_j \rangle \in C_{ij}$
 - ∀ denotes "for all," ∃ denotes "there exists" and ∈ denotes "in."

Revise: A directed arc consistency procedure

Definition: arc $\langle x_i, x_j \rangle$ is directed arc consistent if $\forall a_i \in D_i$, $\exists a_j \in D_j$ such that $\langle a_i, a_j \rangle \in C_{ij}$.

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Revise (x_i, x_j)
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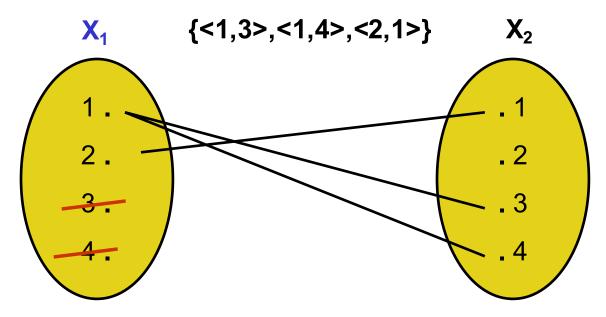
Input: Variables x_i and x_j with domains D_i and D_j and constraint relation R_{ij} . **Output:** pruned D_i , such that x_i is directed arc-consistent relative to x_i .

- 1. for each $a_i \in D_i$
- 2. **if** there is **no** $a_i \in D_i$ such that $\langle a_i, a_i \rangle \in R_{ii}$,
- 3. then delete a_i from D_i.
- 4. endif
- 5. endfor

Constraint Processing, by R. Dechter pgs 54-56.

Directed Arc Consistency

Revise (x_1, x_2) :



Now arc $\langle x_1, x_2 \rangle$ is directed arc consistent.

Definition: arc $\langle x_i, x_j \rangle$ is arc consistent if $\langle x_i, x_j \rangle$ and $\langle x_j, x_i \rangle$ are directed arc consistent.

Definition: Problem is arc consistent if all pairs of variables are arc consistent.

Full Arc Consistency over All Constraints via Constraint Propagation

Definition: arc $\langle x_i, x_j \rangle$ is directed arc consistent if

 \forall $a_i \in D_i$, \exists $a_j \in D_j$ such that $\langle a_i, a_j \rangle \in C_{ij}$.

Constraint Propagation:

To achieve (directed) arc consistency over CSP:

- 1. For every arc C_{ii} in CSP, with tail domain D_i, call Revise.
- 2. Repeat until quiescence:

If an element was deleted from D_i, then

repeat Step 1.

(AC-1)

Full Arc-Consistency via AC-1

AC-1(CSP)

Input: A constraint satisfaction problem CSP = <X, D, C>.

Output: CSP', the largest arc-consistent subset of CSP.

- 1. repeat
- for every c_{ii} ∈ C,
- 3. Revise(x_i , x_j)
- 4. Revise(x_i , x_i)
- 5. **endfor**
- 6. until no domain is changed.

```
For every arc, prune head and tail domains.
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Constraint Processing,
by R. Dechter
pgs 57.
```

Full Arc Consistency via Constraint Propagation

Definition: arc $\langle x_i, x_i \rangle$ is directed arc consistent if

 \forall $a_i \in D_i$, \exists $a_j \in D_j$ such that $\langle a_i, a_j \rangle \in C_{ij}$.

Constraint Propagation:

To achieve (directed) arc consistency over CSP:

- 1. For every arc C_{ii} in CSP, with tail domain D_i, call Revise.
- 2. Repeat until quiescence:

If an element was deleted from D_i, then

```
repeat Step 1,

OR call Revise on each arc with head D<sub>i</sub>

(use FIFO Q, and remove duplicates).

(AC-1)
```

Full Arc-Consistency via AC-3 (Waltz CP)

AC-3(CSP)

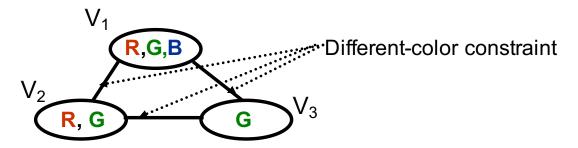
Input: A constraint satisfaction problem CSP = <X, D, C>.

Output: CSP', the largest arc-consistent subset of CSP.

```
Constraint Processing,
1. for every c_{ii} \in C,
                                                                                    by R. Dechter
         queue \leftarrow queue \cup \{\langle x_i, x_i \rangle, \langle x_i, x_i \rangle\}
                                                                                      pgs 58-59.
     endfor
4. while queue ≠ {}
5.
        select and delete arc <x<sub>i</sub>, x<sub>i</sub>> from queue
           Revise(x_i, x_i)
6.
7.
           if Revise(x<sub>i</sub>, x<sub>i</sub>) caused a change in D<sub>i</sub>
              then queue \leftarrow queue \cup \{\langle x_k, x_i \rangle \mid k \neq i, k \neq j\}
8.
9.
           endif
                                          16.410/16.413 Fall 2015
10. endwhile
```

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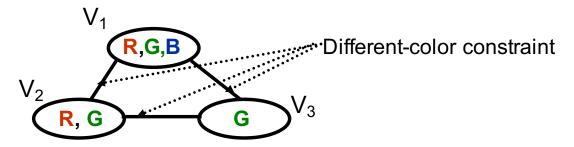
Graph Coloring
Initial Domains



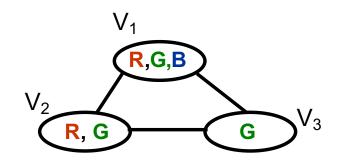
Each undirected arc denotes two directed arcs.

Graph Coloring

Initial Domains



Arc examined	Value deleted

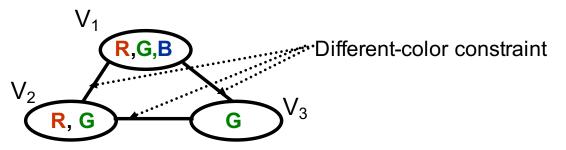


$$V_1 - V_2, V_1 - V_3, V_2 - V_3$$

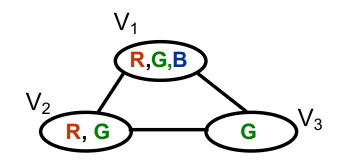
- Introduce queue of arcs to be examined.
- Start by adding all arcs to the queue.

Graph Coloring

Initial Domains



Arc examined	Value deleted

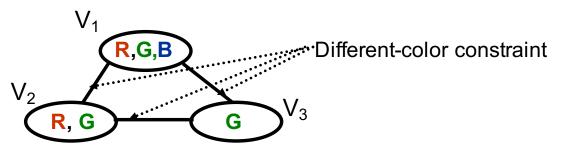


$$V_1 - V_2, V_1 - V_3, V_2 - V_3$$

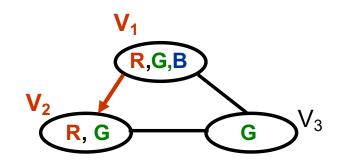
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 > V_2$	

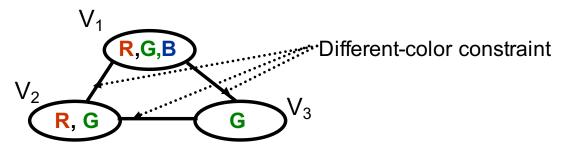


$$V_2 > V_1, V_1 - V_3, V_2 - V_3$$

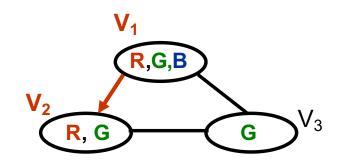
- Delete disallowed tail values.
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 > V_2$	none

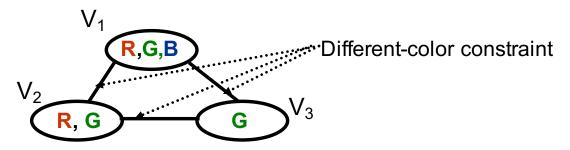


$$V_2 > V_1, V_1 - V_3, V_2 - V_3$$

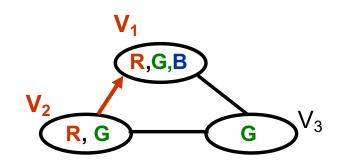
- Delete disallowed tail values.
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 > V_2$	none
$V_2 > V_1$	

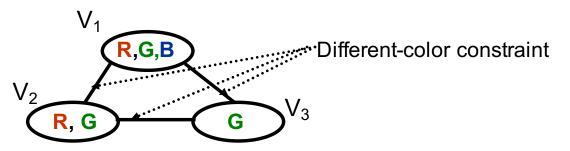


$$V_1 - V_3, V_2 - V_3$$

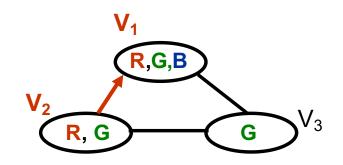
- Delete disallowed tail values.
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 > V_2$	none
$V_2 > V_1$	none

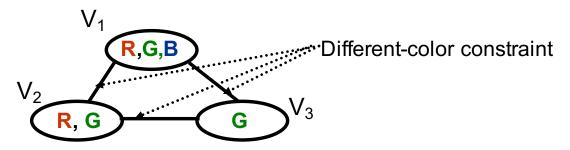


$$V_1 - V_3$$
, $V_2 - V_3$

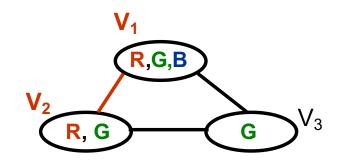
- Delete disallowed tail values.
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none

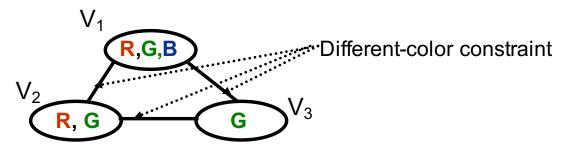


$$V_1 - V_3, V_2 - V_3$$

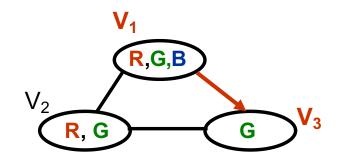
- Delete disallowed tail values.
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 > V_3$	

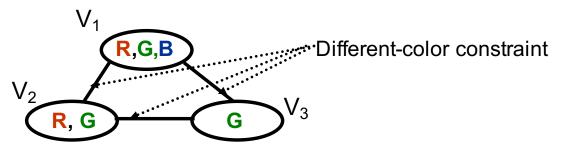


$$V_3 > V_1, V_2 - V_3$$

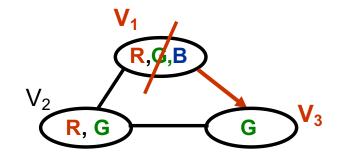
- Delete disallowed tail values.
- $V_i V_j$ denotes two arcs, between V_i and V_j .
- $V_i > V_j$ denotes an arc from V_i to V_j .

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 > V_3$	V ₁ (G)



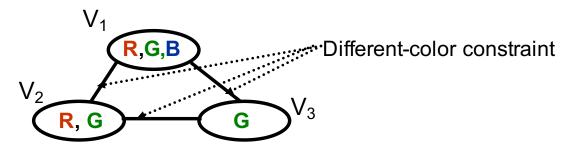
Arcs to examine

$$V_3 > V_1, V_2 - V_3$$

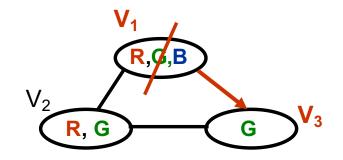
IF THEN An element of a variable's domain is removed, add all arcs to that variable to the examination queue.

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 > V_3$	V ₁ (G)



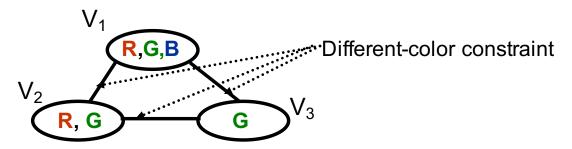
Arcs to examine

$$V_3 > V_1, V_2 - V_3, V_2 > V_1, V_3 > V_1$$

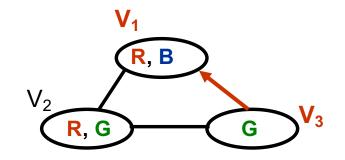
IF THEN An element of a variable's domain is removed, add all arcs to that variable to the examination queue.

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 > V_3$	V ₁ (G)
$V_3 > V_1$	



Arcs to examine

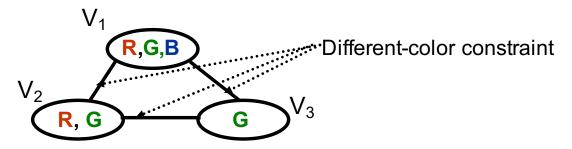
$$V_2 - V_3, V_2 > V_1$$

Delete unmentioned tail values.

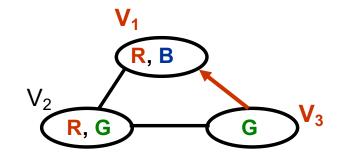
IF THEN An element of a variable's domain is removed, add all arcs to that variable to the examination queue.

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 > V_3$	V ₁ (G)
$V_3 > V_1$	none



Arcs to examine

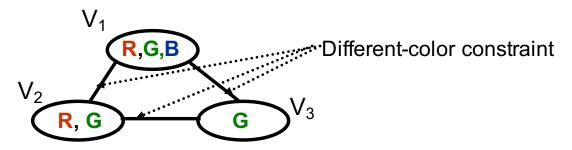
$$V_2 - V_3, V_2 > V_1$$

Delete unmentioned tail values.

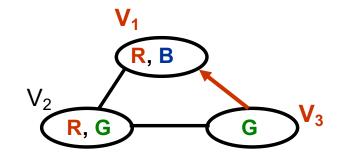
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)



Arcs to examine

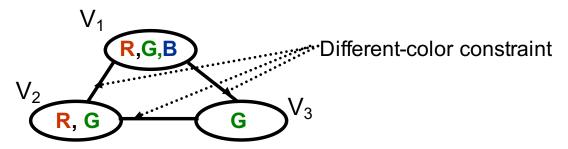
$$V_2 - V_3, V_2 > V_1$$

Delete unmentioned tail values.

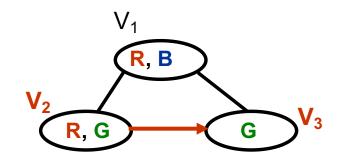
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 > V_3$	



Arcs to examine

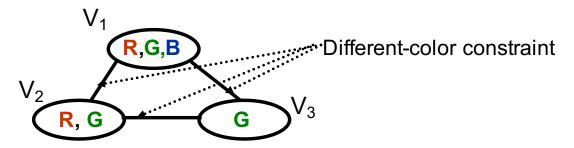
 $V_3 > V_2, V_2 > V_1$

Delete unmentioned tail values.

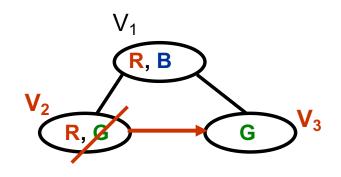
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 > V_3$	V ₂ (G)



Arcs to examine

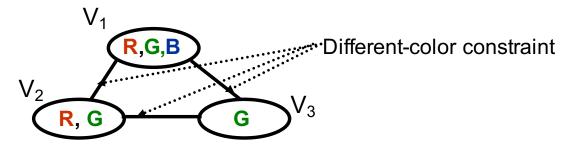
 $V_3 > V_3, V_2 > V_1$

Delete unmentioned tail values.

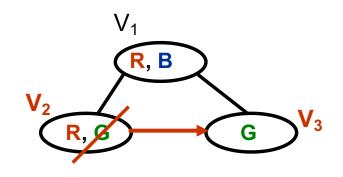
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 > V_3$	V ₂ (G)



Arcs to examine

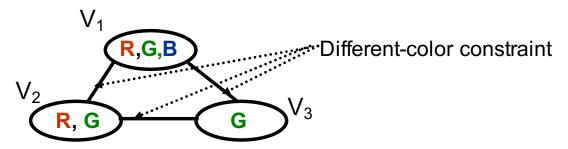
$$V_3 > V_2, V_2 > V_1, V_1 > V_2$$

Delete unmentioned tail values.

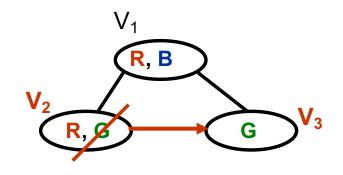
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 > V_3$	V ₂ (G)



Arcs to examine

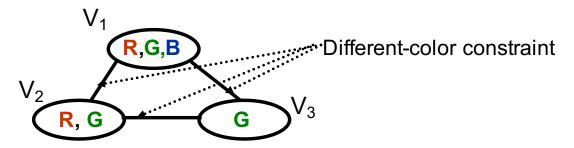
$$V_3 > V_2, V_2 > V_1, V_1 > V_2$$

Delete unmentioned tail values.

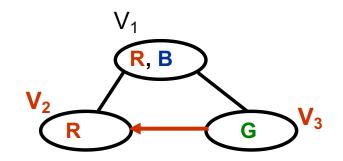
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 > V_3$	V ₂ (G)
$V_3 > V_2$	



Arcs to examine

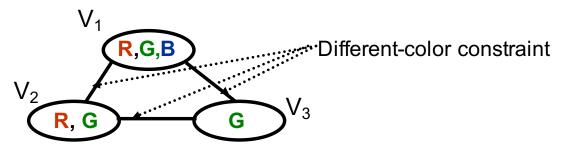
 $V_2 > V_1, V_1 > V_2$

Delete unmentioned tail values.

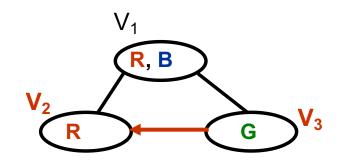
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 > V_3$	V ₂ (G)
$V_3 > V_2$	none



Arcs to examine

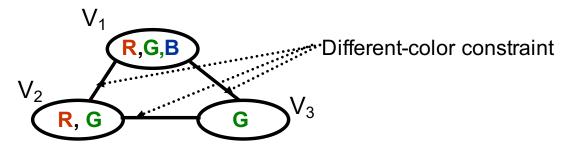
$$V_2 > V_1, V_1 > V_2$$

Delete unmentioned tail values.

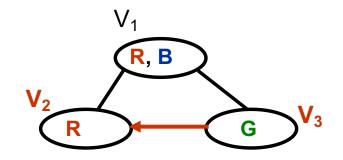
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)



Arcs to examine

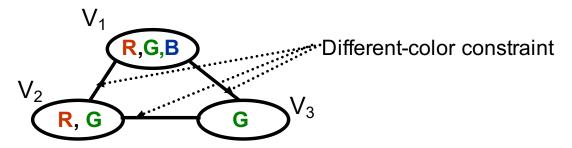
$$V_2 > V_1, V_1 > V_2$$

Delete unmentioned tail values.

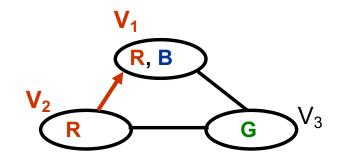
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
V ₂ > V ₁	



Arcs to examine

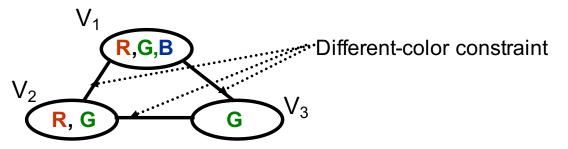
 $V_1 > V_2$

Delete unmentioned tail values.

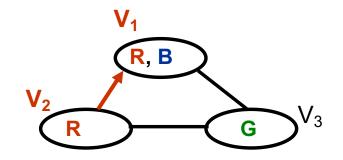
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 > V_1$	none



Arcs to examine

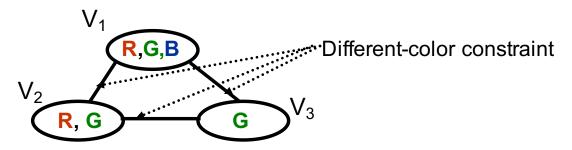
 $V_1 > V_2$

Delete unmentioned tail values.

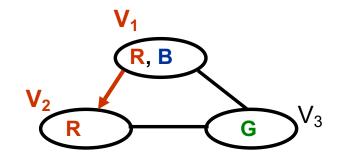
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 > V_1$	none
$V_1 > V_2$	



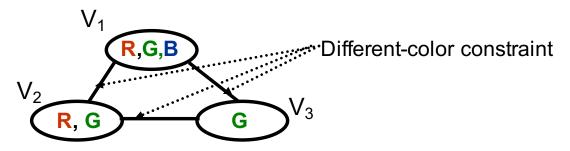
Arcs to examine

Delete unmentioned tail values.

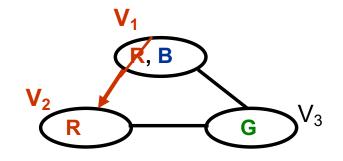
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 > V_1$	none
$V_1 > V_2$	V ₁ (R)



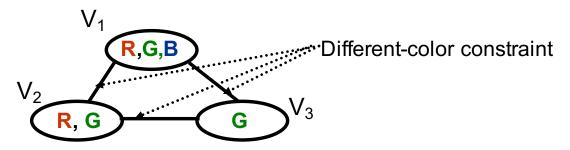
Arcs to examine

Delete unmentioned tail values.

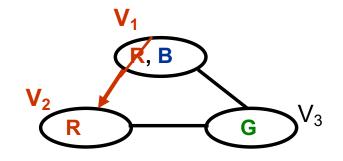
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 > V_1$	none
$V_1 > V_2$	V ₁ (R)



Arcs to examine

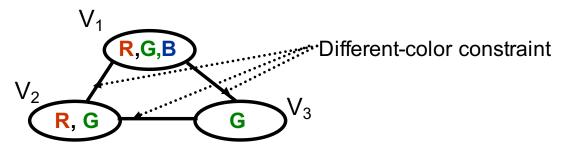
$$V_2 > V_1, V_3 > V_1$$

Delete unmentioned tail values.

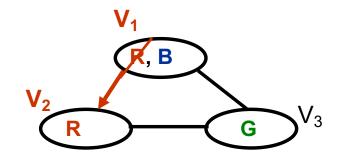
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 - V_1$	V ₁ (R)



Arcs to examine

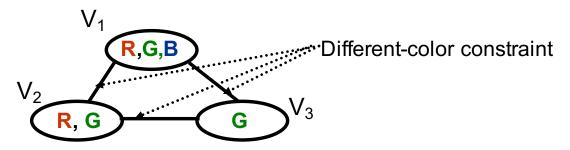
$$V_2 > V_1, V_3 > V_1$$

Delete unmentioned tail values.

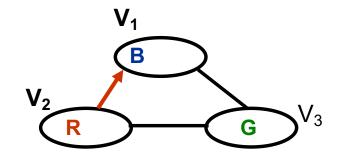
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 - V_1$	V ₁ (R)
$V_2 > V_1$	



Arcs to examine

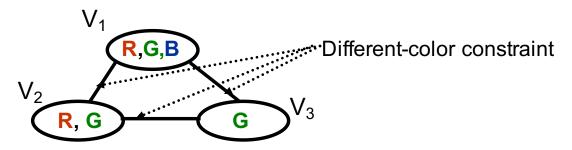
 $V_3 > V_1$

Delete unmentioned tail values.

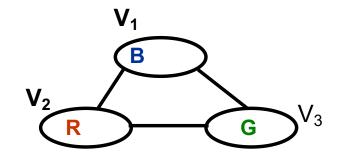
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 - V_1$	V ₁ (R)
$V_2 > V_1$	none



Arcs to examine

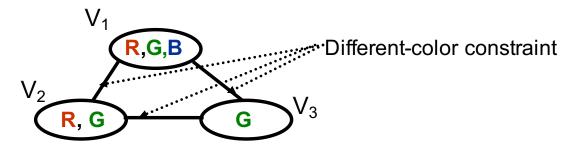
 $V_3 > V_1$

Delete unmentioned tail values.

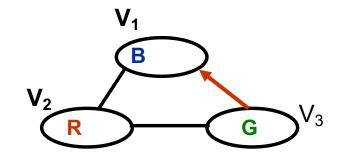
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 - V_1$	V ₁ (R)
$V_2 > V_1$	none
$V_3 > V_1$	



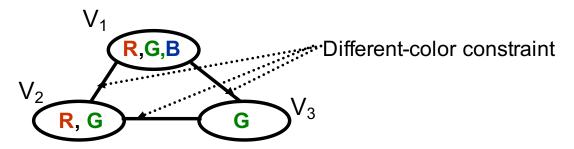
Arcs to examine

Delete unmentioned tail values.

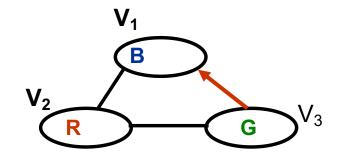
IF THEN

Graph Coloring

Initial Domains



Arc examined	Value deleted
$V_1 - V_2$	none
$V_1 - V_3$	V ₁ (G)
$V_2 - V_3$	V ₂ (G)
$V_2 - V_1$	V ₁ (R)
$V_2 > V_1$	none
$V_3 > V_1$	none



Arcs to examine

IF examination queue is empty

THEN arc (pairwise) consistent.

Outline

- Arc-consistency and constraint propagation.
- Analysis of constraint propagation.
- Solving CSPs using search.

What is the Complexity of AC-1?

AC-1(CSP)

Input: A network of constraints CSP = <X, D, C>.

Output: CSP', the largest arc-consistent subset of CSP.

- 1. repeat
- **2. for** every $c_{ij} \in C$,
- 3. Revise(x_i , x_i)
- 4. Revise(x_i , x_i)
- 5. **endfor**
- 6. until no domain is changed.

Assume:

- There are n variables.
- Domains are of size at most k.
- There are e binary constraints.

What is the Complexity of AC-1?

Assume:

- There are n variables.
- Domains are of size at most k.
- There are e binary constraints.

Which is the correct complexity?

- 1. $O(k^2)$,
- 2. O(enk²),
- 3. $O(enk^3)$,
- 4. O(nek).

Revise: A directed arc consistency procedure

Revise (x_i, x_j)

Input: Variables x_i and x_i with domains D_i and D_i and constraint relation R_{ii} .

Output: pruned D_i , such that x_i is directed arc-consistent relative to x_i .

- 1. for each $a_i \in D_i$
- 2. **if** there is no $a_i \in D_i$ such that $\langle a_i, a_i \rangle \in R_{ii}$
- 3. then delete a_i from D_i .
- 4. endif
- 5. endfor

Complexity of Revise? $= O(k^2)$.

where $k = \max_{i} |D_{i}|$

O(k)

* O(k)

Full Arc-Consistency via AC-1

```
AC-1(CSP)
Input: A network of constraints CSP = <X, D, C>.
Output: CSP', the largest arc-consistent subset of CSP.
    repeat
                                                                O(2e*revise)
     for every c_{ij} \in C,
       Revise(x_i, x_i)
3.
       Revise(x_i, x_i)
4.
5.
      endfor
                                                                * O(nk)
    until no domain is changed.
Complexity of AC-1?
        = O(nk*e*revise),
```

 $= O(enk^3),$

where
$$k = max_i |D_i|$$
,
16.410/16.413 Fall 2015 $n = |X|$, $e = |C|$.

What is the Complexity of Constraint Propagation using AC-3?

Assume:

- There are n variables.
- Domains are of size at most k.
- There are e binary constraints.

Which is the correct complexity?

- 1. $O(k^2)$,
- 2. $O(ek^2)$,
- 3. $O(ek^3)$,
- 4. O(ek).

Full Arc-Consistency via AC-3

```
AC-3(CSP)
Input: A network of constraints CSP = <X, D, C>.
Output: CSP', the largest arc-consistent subset of CSP.
     for every c_{ii} \in C,
                                                                                  O(e) +
        queue \leftarrow queue \cup \{\langle x_i, x_i \rangle, \langle x_i, x_i \rangle\}
3.
     endfor
     while queue ≠ {}
5.
       select and delete arc <x<sub>i</sub>, x<sub>i</sub>> from queue
                                                                                   O(k^2)
6.
         Revise(x_i, x_i)
                                                                                    * O(ek)
7.
         if Revise(x_1, x_2) caused a change in D_i.
            then queue \leftarrow queue \cup \{\langle x_k, x_l \rangle \mid k \neq i, k \neq j\}
8.
         endif
9_
10. endwhile
Complexity of AC-3?
                                                    where k = max_{\cdot} |D_{\cdot}|, n = |X|, e = |C|.
          = O(e+ek*k^2) = O(ek^3),
```

16.410/16.413 Fall 2015

Is arc consistency sound and complete?

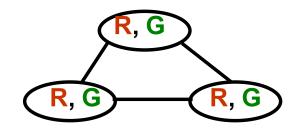
An arc consistent solution selects a value for every variable from its arc consistent domain.

Soundness: All solutions to the CSP are arc consistent solutions?

- •Yes,
- · No.

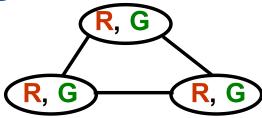
Completeness: All arc-consistent solutions are solutions to the CSP?

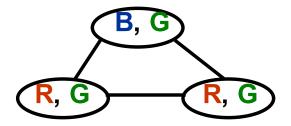
- Yes,
- · No.



Incomplete: Arc consistency doesn't rule out all infeasible solutions

Graph Coloring





Arc consistent, but no solutions.

Arc consistent, but <u>2</u> solutions, not 8.

To Solve CSPs We Combine

1. Arc consistency (via constraint propagation):

Eliminates values that are shown locally to not be a part of any solution.

2. Search:

Explores consequences of committing to particular assignments.

Methods that Incorporate Search:

- Standard Search,
- Back Track Search (BT),
- BT with Forward Checking (FC),
- Dynamic Variable Ordering (DV),
- Iterative Repair (IR),
- Conflict-directed Back Jumping (CBJ).

Solving CSPs using Generic Search

State

 Partial assignment to variables, made thus far.

Initial State

· No assignment.

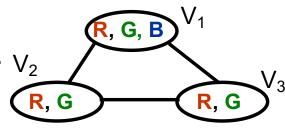
Operator

- Creates new assignment = (X_i = v_{ii}).
 - Select any unassigned variable X_i.
 - Select any one of its domain values v_{ii}.
- · Child extends parent assignments with new.

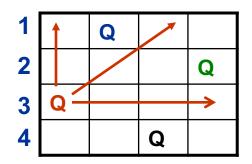
Goal Test

- All variables are assigned.
- All constraints are satisfied.

- Branching factor?
- \rightarrow Sum of domain size of all variables $O(|v|^*|d|)$. V_2
- Performance?
- \rightarrow Exponential in the branching factor $O([|v|^*|d|]^{|v|})$.



Search Performance on N Queens



Standard Search,

A handful of queens.

Backtracking.

Solving CSPs with Standard Search

Standard Search:

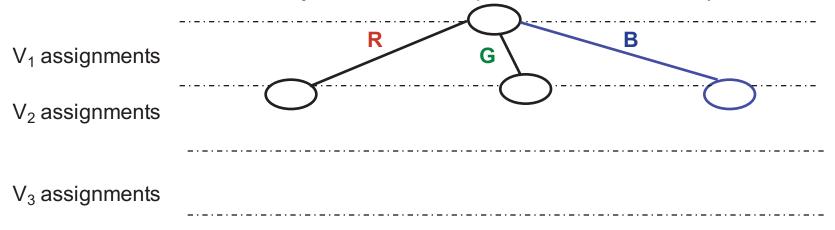
- Children select any value for any variable [O(|v|*|d|)].
- Test complete assignments for consistency against CSP.

Observations:

- 1. The order in which variables are assigned does not change the solution.
 - Many paths denote the same solution,
 - (|v|!).
 - → Expand only one path (i.e., use one variable ordering).
- 1. We can identify and prune a dead end before we assign all variables.
 - Extensions to inconsistent partial assignments are always inconsistent.
 - Check consistency after each assignment.

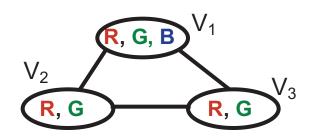
Next: Back Track Search (BT)

- 1. Expand assignments of one variable at each step.
- 2. Pursue depth first.
- 3. Check consistency after each expansion, and backup.



Preselect order of variables to assign.

Assign designated variable.



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