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6.080 / 6.089 Great Ideas in Theoretical Computer Science  
Spring 2008

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## 6.080/6.089 Problem Set 5

Assigned: April 29, 2008

Due: May 13, 2008

1. Let a *puzzle generator* be a polynomial-time algorithm that maps a random string  $r$  to a pair  $(\varphi_r, x_r)$ , where  $\varphi_r$  is a 3SAT instance and  $x_r$  is a satisfying assignment for  $\varphi_r$ , such that for all polynomial-time algorithms  $A$ ,

$$\Pr_r [A \text{ finds a satisfying assignment for } \varphi_r]$$

is negligible (less than  $\frac{1}{\text{poly}(n)}$ ). Show that puzzle generators exist if and only if one-way functions exist.

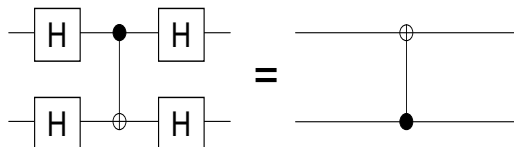
2. The following questions concern the RSA cryptosystem.

- (a) Recall that, having chosen primes  $p$  and  $q$  such that  $p - 1$  and  $q - 1$  are not divisible by 3, a key step in RSA is to find an integer  $k$  such that  $3k \equiv 1 \pmod{(p - 1)(q - 1)}$ . Give a simple procedure to find such a  $k$  given  $p$  and  $q$ , which requires only  $O(1)$  arithmetic operations.
- (b) Given a product of two primes,  $N = pq$ , show that if an eavesdropper can efficiently determine  $(p - 1)(q - 1)$  (the order of the multiplicative group mod  $N$ ), then she can also efficiently determine  $p$  and  $q$  themselves.

3. Recall that the *VC-dimension* of a concept class  $\mathcal{C}$ , or  $\text{VCdim}(\mathcal{C})$ , is the largest  $m$  for which there exist points  $x_1, \dots, x_m$  such that for all  $2^m$  possible Boolean values of  $c(x_1), \dots, c(x_m)$ , there exists a concept  $c \in \mathcal{C}$  that realizes those values. If such  $x_1, \dots, x_m$  exist for arbitrarily large finite  $m$ , then  $\text{VCdim}(\mathcal{C}) = \infty$ .

- (a) Let  $\mathcal{C}$  be the concept class consisting of all filled-in rectangles in the plane, whose sides are aligned with the  $x$  and  $y$  axes. Show that  $\text{VCdim}(\mathcal{C}) = 4$ .
- (b) Show that if  $\mathcal{C}$  is finite, then  $\text{VCdim}(\mathcal{C}) \leq \log_2 |\mathcal{C}|$ .
- (c) Show that there is a class  $\mathcal{C}$  with countably many concepts such that  $\text{VCdim}(\mathcal{C}) = \infty$ .

4. Show that if you apply Hadamard gates to qubits  $A$  and  $B$ , followed by a CNOT gate from  $A$  to  $B$ , followed by Hadamard gates to  $A$  and  $B$  again, the end result is the same as if you had applied a CNOT gate from  $B$  to  $A$ . Pictorially:



This illustrates a principle of quantum mechanics you may have heard about: that any physical interaction by which  $A$  influences  $B$  can also cause  $B$  to influence  $A$  (so for example, it is impossible to measure a particle's state without affecting it).

5. Consider the following game played by Alice and Bob. Alice receives a bit  $x$  and Bob receives a bit  $y$ , with both bits uniformly random and independent. The players win if Alice outputs a bit  $a$  and Bob outputs a bit  $b$  such that  $a + b = xy \pmod{2}$ . (Alice and Bob are cooperating in this game, not competing.) The players can agree on a strategy in advance, but once they receive  $x$  and  $y$  no further communication between them is allowed.
- (a) Give a deterministic strategy by which Alice and Bob can win this game with  $3/4$  probability.
  - (b) Show that no deterministic strategy lets them win with more than  $3/4$  probability.
  - (c) [*Extra credit*] Show that no probabilistic strategy lets them win with more than  $3/4$  probability.

Now suppose Alice and Bob share the entangled state  $\frac{1}{\sqrt{2}}(|00\rangle + |11\rangle)$ , with Alice holding one qubit and Bob holding the other qubit. Suppose they use the following strategy: if  $x = 1$ , then Alice applies the unitary matrix

$$\begin{pmatrix} \cos \frac{\pi}{8} & -\sin \frac{\pi}{8} \\ \sin \frac{\pi}{8} & \cos \frac{\pi}{8} \end{pmatrix}$$

to her qubit, otherwise she doesn't. She then measures her qubit in the standard basis and outputs the result. If  $y = 1$ , then Bob applies the unitary matrix

$$\begin{pmatrix} \cos \frac{\pi}{8} & \sin \frac{\pi}{8} \\ -\sin \frac{\pi}{8} & \cos \frac{\pi}{8} \end{pmatrix}$$

to his qubit, otherwise he doesn't. He then measures his qubit in the standard basis and outputs the result.

- d. Show that if  $x = y = 0$ , then Alice and Bob win the game with probability 1 using this strategy.
- e. Show that if  $x = 1$  and  $y = 0$  (or vice versa), then Alice and Bob win with probability  $\cos^2 \frac{\pi}{8} = \frac{1 + \sqrt{1/2}}{2}$ .
- f. Show that if  $x = y = 1$ , then Alice and Bob win with probability  $1/2$ .
- g. Combining parts d-f, conclude that Alice and Bob win with greater overall probability than would be possible in a classical universe.

You have just proved the *CHSH/Bell Inequality*—one of the most famous results of quantum mechanics—which showed the impossibility of Einstein's dream of removing "spooky action at a distance" from quantum mechanics. Alice and Bob's ability to win the above game more than  $3/4$  of the time using quantum entanglement was experimentally confirmed in the 1980's.