

Question 1

I first prove that if $\Gamma \mathbf{P} \vee \mathbf{Q} \neg \in \Gamma^*$, then $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$. I do this by proving the contrapositive — i.e., that if it is not the case that $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$, then it is not the case that $\Gamma \mathbf{P} \vee \mathbf{Q} \neg \in \Gamma^*$.

1. Suppose it is not the case that $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$.
2. Then $\Gamma \sim \mathbf{P} \neg \in \Gamma^*$ and $\Gamma \sim \mathbf{Q} \neg \in \Gamma^*$ (by 6.4.11(a)).
3. So $\{\Gamma \sim \mathbf{P} \neg, \Gamma \sim \mathbf{Q} \neg\} \subset \Gamma^*$.
4. Now, $\{\Gamma \sim \mathbf{P} \neg, \Gamma \sim \mathbf{Q} \neg\} \vdash \Gamma \sim (\mathbf{P} \vee \mathbf{Q}) \neg$, in *SD* (proof below).
5. So $\Gamma \sim (\mathbf{P} \vee \mathbf{Q}) \neg \in \Gamma^*$ (by 3, 4 and 6.4.9).
6. So it is not the case that $\Gamma \mathbf{P} \vee \mathbf{Q} \neg \in \Gamma^*$ (again by 6.4.11(a)).

So, if $\Gamma \mathbf{P} \vee \mathbf{Q} \neg \in \Gamma^*$, then $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$. Q.E.D.

Here is a proof of 4 (I think we did something very like this in class, but I do a derivation here anyway for completeness' sake).

1	$\sim \mathbf{P}$	A
2	$\sim \mathbf{Q}$	A
3	$\mathbf{P} \vee \mathbf{Q}$	A/ \sim I
4	\mathbf{P}	A/ \vee E
5	\mathbf{P}	4, R
6	\mathbf{Q}	A/ \vee E
7	$\sim \mathbf{P}$	A/ \sim E
8	\mathbf{Q}	6, R
9	$\sim \mathbf{Q}$	2, R
10	\mathbf{P}	7-9, \sim E
11	\mathbf{P}	3, 4-5, 6-10, \vee E
12	$\sim \mathbf{P}$	1, R
13	$\sim (\mathbf{P} \vee \mathbf{Q})$	3-12, \sim I

Now to prove the other direction: if $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$, then $\Gamma \mathbf{P} \vee \mathbf{Q} \neg \in \Gamma^*$. Again, I do this by proving the contrapositive.

1. Suppose it is not the case that $\Gamma \mathbf{P} \vee \mathbf{Q} \neg \in \Gamma^*$.
2. Then $\Gamma \sim (\mathbf{P} \vee \mathbf{Q}) \neg \in \Gamma^*$ (by 6.4.11(a)).
3. So $\{\Gamma \sim (\mathbf{P} \vee \mathbf{Q}) \neg\} \subset \Gamma^*$.

4. Now, $\{\ulcorner \sim (\mathbf{P} \vee \mathbf{Q}) \urcorner\} \vdash \ulcorner \sim \mathbf{P} \urcorner$ in SD , and $\{\ulcorner \sim (\mathbf{P} \vee \mathbf{Q}) \urcorner\} \vdash \ulcorner \sim \mathbf{Q} \urcorner$ in SD (proof below).
5. So, $\ulcorner \sim \mathbf{P} \urcorner \in \Gamma^*$ and $\ulcorner \sim \mathbf{Q} \urcorner \in \Gamma^*$ (by 3, 4 and 6.4.9).
6. So it is not the case that either $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$ (by 6.4.11(a) again).

So, if $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$, then $\ulcorner \mathbf{P} \vee \mathbf{Q} \urcorner \in \Gamma^*$. Q.E.D.

Here is a proof of the first half of 4 — i.e., that $\{\ulcorner \sim (\mathbf{P} \vee \mathbf{Q}) \urcorner\} \vdash \ulcorner \sim \mathbf{P} \urcorner$.

1	$\sim (\mathbf{P} \vee \mathbf{Q})$	A						
2	<table style="border-collapse: collapse; border-left: 1px solid black;"> <tr> <td style="padding-right: 10px;">2</td> <td style="padding-left: 10px;">\mathbf{P}</td> <td style="padding-left: 10px;">A/\simI</td> </tr> <tr> <td style="padding-right: 10px;">3</td> <td style="padding-left: 10px;">$\mathbf{P} \vee \mathbf{Q}$</td> <td style="padding-left: 10px;">2, \veeI</td> </tr> </table>	2	\mathbf{P}	A/ \sim I	3	$\mathbf{P} \vee \mathbf{Q}$	2, \vee I	
2	\mathbf{P}	A/ \sim I						
3	$\mathbf{P} \vee \mathbf{Q}$	2, \vee I						
4	$\sim (\mathbf{P} \vee \mathbf{Q})$	1, R						
5	$\sim \mathbf{P}$	2-4, \sim I						

The proof of the other half of 4 is the same, except you replace the ‘ \mathbf{P} ’s on lines 2 and 5 with ‘ \mathbf{Q} ’s.

So, I’ve proven that if $\ulcorner \mathbf{P} \vee \mathbf{Q} \urcorner \in \Gamma^*$, then $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$, and I’ve proven that if $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$, then $\ulcorner \mathbf{P} \vee \mathbf{Q} \urcorner \in \Gamma^*$. It follows that $\ulcorner \mathbf{P} \vee \mathbf{Q} \urcorner \in \Gamma^*$ if and only if $\mathbf{P} \in \Gamma^*$ or $\mathbf{Q} \in \Gamma^*$. And that concludes the proof.

Question 2

We’re trying to prove *Inductive Step*, on p. 273 of TLB, for the case in which \mathbf{P} , a sentence containing $\mathbf{k} + 1$ occurrences of connectives, has the form $\ulcorner \mathbf{Q} \vee \mathbf{R} \urcorner$.

1. Suppose that every sentence of SL with \mathbf{k} or fewer occurrences of connectives is such that it is true on \mathbf{A}^* if and only if it is a member of Γ^* (i.e., suppose the antecedent of *Inductive Step*).
2. Now, $\ulcorner \mathbf{Q} \vee \mathbf{R} \urcorner$ is true on \mathbf{A}^* iff either \mathbf{Q} is true on \mathbf{A}^* or \mathbf{R} is true on \mathbf{A}^* (by definition of ‘ \vee ’).
3. And \mathbf{Q} is true on \mathbf{A}^* iff $\mathbf{Q} \in \Gamma^*$, and \mathbf{R} is true on \mathbf{A}^* iff $\mathbf{R} \in \Gamma^*$ (by 1, and the fact that \mathbf{Q}, \mathbf{R} both contain \mathbf{k} or fewer occurrences of connectives).
4. So $\ulcorner \mathbf{Q} \vee \mathbf{R} \urcorner$ is true on \mathbf{A}^* iff either $\mathbf{Q} \in \Gamma^*$ or $\mathbf{R} \in \Gamma^*$ (from 2, 3).
5. So $\ulcorner \mathbf{Q} \vee \mathbf{R} \urcorner$ is true on \mathbf{A}^* if and only if $\ulcorner \mathbf{Q} \vee \mathbf{R} \urcorner \in \Gamma^*$ (by 6.4.11(c) — i.e., the thing we just proved in Question 1).

So *Inductive Step* is true for the case in which \mathbf{P} has the form $\ulcorner \mathbf{Q} \vee \mathbf{R} \urcorner$. Q.E.D.

Question 3

The completeness proof for SD will fail, as a proof for the completeness of SD^* , at the part where we try to prove **6.4.11(b)** — i.e., the proof that $\lceil \mathbf{P} \& \mathbf{Q} \rceil \in \Gamma^*$ if and only if both $\mathbf{P} \in \Gamma^*$ and $\mathbf{Q} \in \Gamma^*$ (where Γ^* is a maximal consistent-in- SD set of sentence of SL ; \mathbf{P} , \mathbf{Q} are sentence of SL) will not go through. In particular, the proof that if $\lceil \mathbf{P} \& \mathbf{Q} \rceil \in \Gamma^*$ then both $\mathbf{P} \in \Gamma^*$ and $\mathbf{Q} \in \Gamma^*$ will not go through. Note that the proof of that part of **6.4.11(b)**, on p. 272 of TLB, involves appealing to the Conjunction Elimination rule explicitly.

In fact, it will not, in general, be the case that a maximal consistent-in- SD^* set is such that if $\lceil \mathbf{P} \& \mathbf{Q} \rceil \in \Gamma^*$ then both $\mathbf{P} \in \Gamma^*$ and $\mathbf{Q} \in \Gamma^*$ (though this is quite hard to prove, and I don't do so here). There will, for example, be maximal consistent-in- SD^* sets that are supersets of $\{ 'A \& B', '\sim A' \}$.

Because the proof of **6.4.11(b)** fails, the proof of what the book calls the 'Consistency Lemma' fails too; in particular, case 2 of the inductive step fails. Even more in particular, the part of case 2 in which we prove that if $\lceil \mathbf{Q} \& \mathbf{R} \rceil$ is false on \mathbf{A}^* than it is not in Γ^* will fail. That part of the proof relies on the part of **6.4.11(b)** that fails without Conjunction Elimination. And you can see why: for a set that is a maximal consistent-in- SD^* superset of $\{ 'A \& B', \sim A' \}$, $'A \& B'$ will be false on \mathbf{A}^* , but it is in there anyway.

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