

mas.s62
lecture 15
discreet log contracts

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Tadge Dryja

today

discreet log contracts

conditional payments

oracles

anticipated signatures

building discreet log contracts

conditional payments

payment conditional on some external data

In this example, Alice and Bob bet on tomorrow's weather. If it rains, Alice gets 1 BTC. If it's sunny, Bob gets 1 BTC.

One problem: The bitcoin blockchain is not aware of the weather. (OP_WEATHER has not yet been soft-forked in)

"smart contracts" and oracles

LN is a simple script, enforcing the most recent tx

Made of smart contracts, but has no external state. Everything comes from Alice & Bob

If we want external state, need some way to get it, usually called an "oracle"

Simple oracle: 2 of 3 multisig

why oracles?

2 of 2 multisig means conflict freezes funds

Rich players at an advantage (lower time value of money)

Works great with friends, but bitcoin is the currency of enemies :)

A 3rd party can decide in case of conflict

2 of 3 multisig oracle

2 of 3 multisig oracle

3 keys: Alice, Bob, Olivia

If Alice and Bob are chill, they can both sign without contacting Olivia

If Alice and Bob fight or are unresponsive, one of them can ask Olivia to sign

Problem: It's sunny. Alice tells Olivia, "Hey, Alice. Say it's raining and I'll give you 0.8"

oracle interaction

2 of 3 multisig oracles are interactive

Not only do they see every contract, they decide the outcome of every contract, individually. (Can equivocate)

It'd be better if the oracle couldn't equivocate, and even better if they never saw the contracts. But how?

revokable tx

Commit Tx (held by Alice)	
input	output
fund txid Bob's signature	Alice key & 100 blocks or AliceR & Bob key 2 coins
	Bob address 8 coins

revokable tx

Commit Tx (held by Bob)	
input	output
fund txid Alice's signature	Alice address 2 coins
	Bob key & 100 blocks or Alice & BobR key 8 coins

point and scalar operations

(Note also works on exponents mod n)

a, b lowercase = scalar

A, B uppercase = point

what operations can we do?

point and scalar operations

scalars are regular unleaded numbers

$a+b$ $a-b$ $a*b$ a/b

everything is OK! just numbers!

point and scalar operations

Points have addition defined.. but not multiplication and division (group)

A+B A-B OK **A*B A/B NO**

add & subtract OK, but can't multiply two points, or divide a point by a point. Not defined.

point and scalar operations

Mixed operations

$A+b$ $A-b$ NO $A*b$ A/b OK

adding points and scalars is undefined

point times scalar OK; repeat the tangent doubling process. Division by scalar also possible.

point and scalar operations

roster of ops: what can we do

$a+b$ $a-b$ $a*b$ a/b (obvious)

$A+B$ $A-B$ $A*b$ A/b

point and scalar operations

roster of ops: what can we do

$a+b$ $a-b$ $a*b$ a/b (obvious)

$A+B$ $A-B$ $A*b$ A/b

Pick some random point G

That's the generator point

Everyone agrees on G

adding pubkeys

$$(aG) + (bG) = (a+b)G$$

sum of private keys gives sum of
public keys! fun stuff ensues

adding pubkeys

$$aG = A, \quad bG = B$$

$$A+B = C = (a+b)G$$

Alice knows a , Bob knows b . Neither can sign with C .

Bob can give b to Alice, then Alice can sign with C .

discreet log contracts
smart contracts in same channel
construction as lightning
lightning: most recent tx is valid
DLC: non-interactive oracle
determines valid tx

schnorr signature

public key $A = aG$

$k \leftarrow \$$; $R = kG$ (nonce for signature)

to sign, compute $s = k - h(m, R)a$

signature is (R, s)

To verify $sG \stackrel{?}{=} kG - h(m, R)aG$

$\stackrel{?}{=} R - h(m, R)A$

fixed-R signature

Pubkey: A signature: (R, s)

Pubkey: (A, R) signature: s

Same thing right? Just move the R.
But can only sign once!

k-collision

$$\text{Signature 1} \quad s_1 = k - h(m_1, R)a$$

$$\text{Signature 2} \quad s_2 = k - h(m_2, R)a$$

$$s_1 - s_2 = k - h(m_1, R)a - k + h(m_2, R)a$$

$$= h(m_2, R)a - h(m_1, R)a$$

$$= (h(m_2, R) - h(m_1, R))a$$

$$a = (s_1 - s_2) / (h(m_2, R) - h(m_1, R))$$

anticipated signature

Given 'pubkey' (A, R) and a message m , you can't compute s .

(EC Discrete log problem)

but you CAN compute $sG = R - h(m, R)A$

sG is computable for any message!

signatures as private keys

It's an unknown scalar, but you know what it is times the generator point.

Hmm! Sounds like a keypair!

Use for a 3rd party oracle to sign messages, revealing a private key.

signatures as private keys

Olivia's s as private key

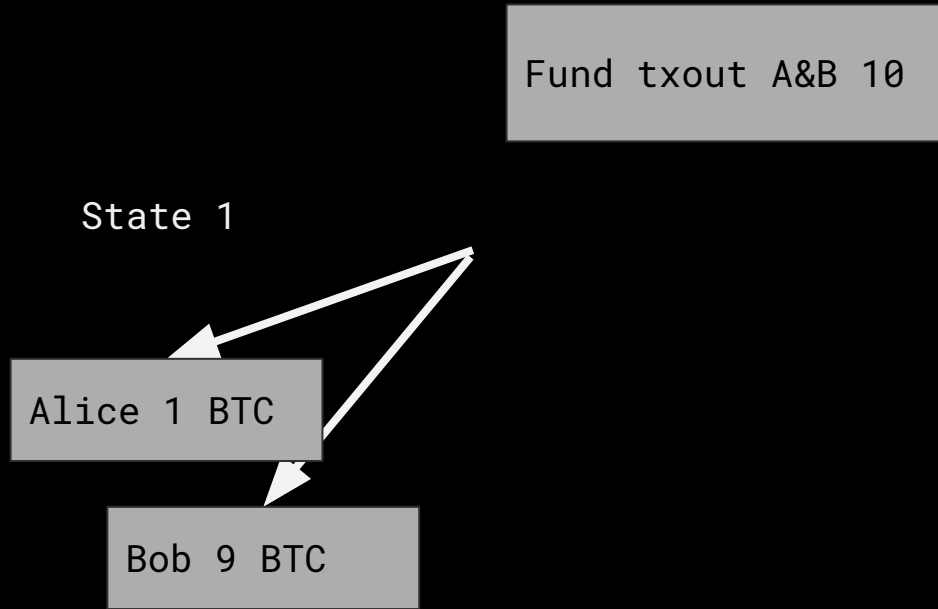
sG as public key

Mix with Alice and Bob's public keys

$$\text{pub}_{\text{alice}} + sG = \text{pub}_{\text{contract}}$$

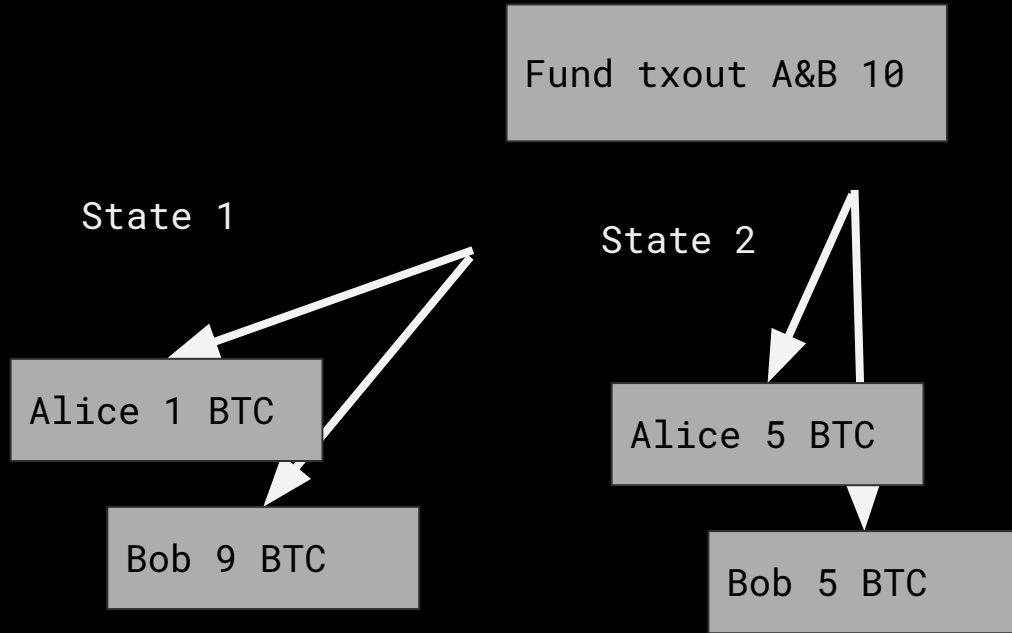
$$\text{priv}_{\text{alice}} + s = \text{priv}_{\text{contract}}$$

signatures as private keys



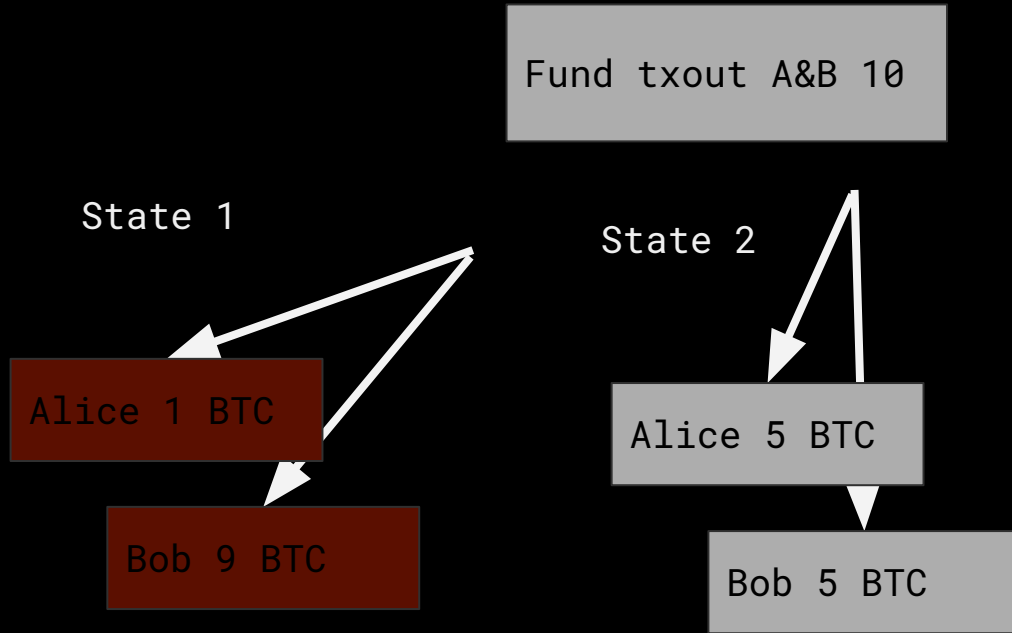
In Lightning, states are added sequentially, and validity is enforced by revealing private keys to previous states

signatures as private keys



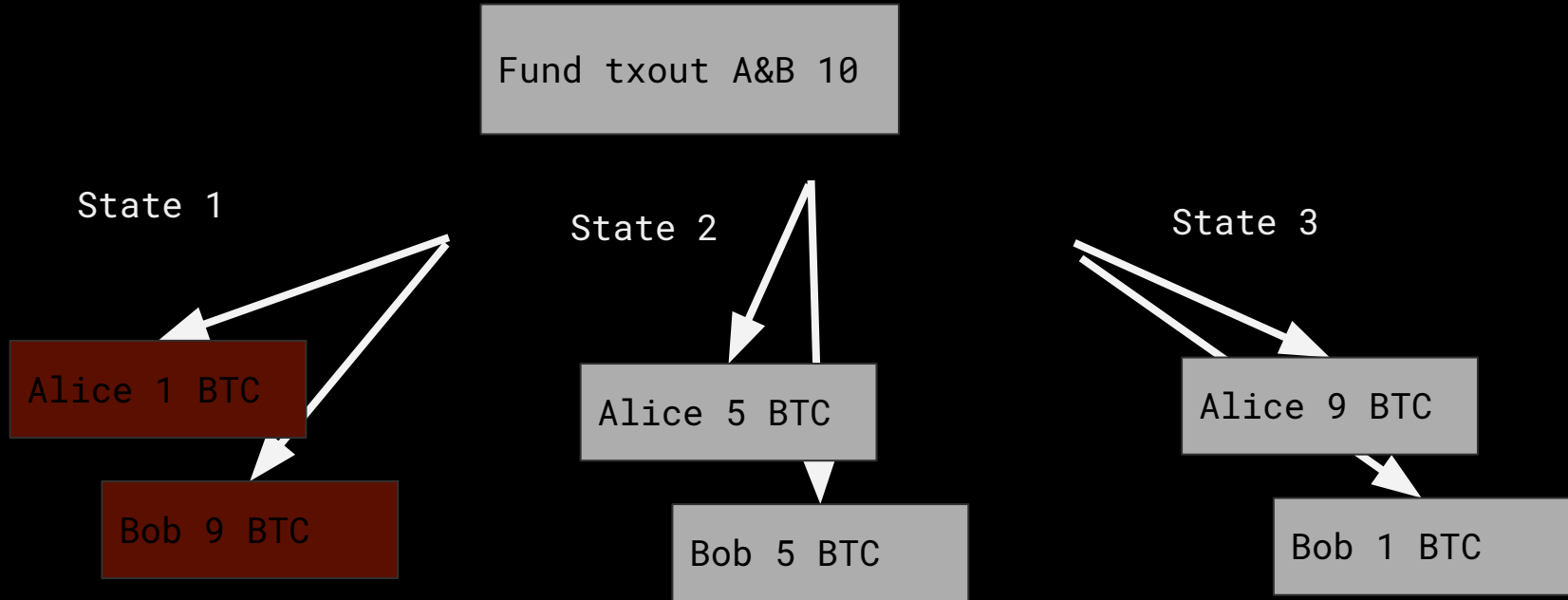
In Lightning, states are added sequentially, and validity is enforced by revealing private keys to previous states

signatures as private keys



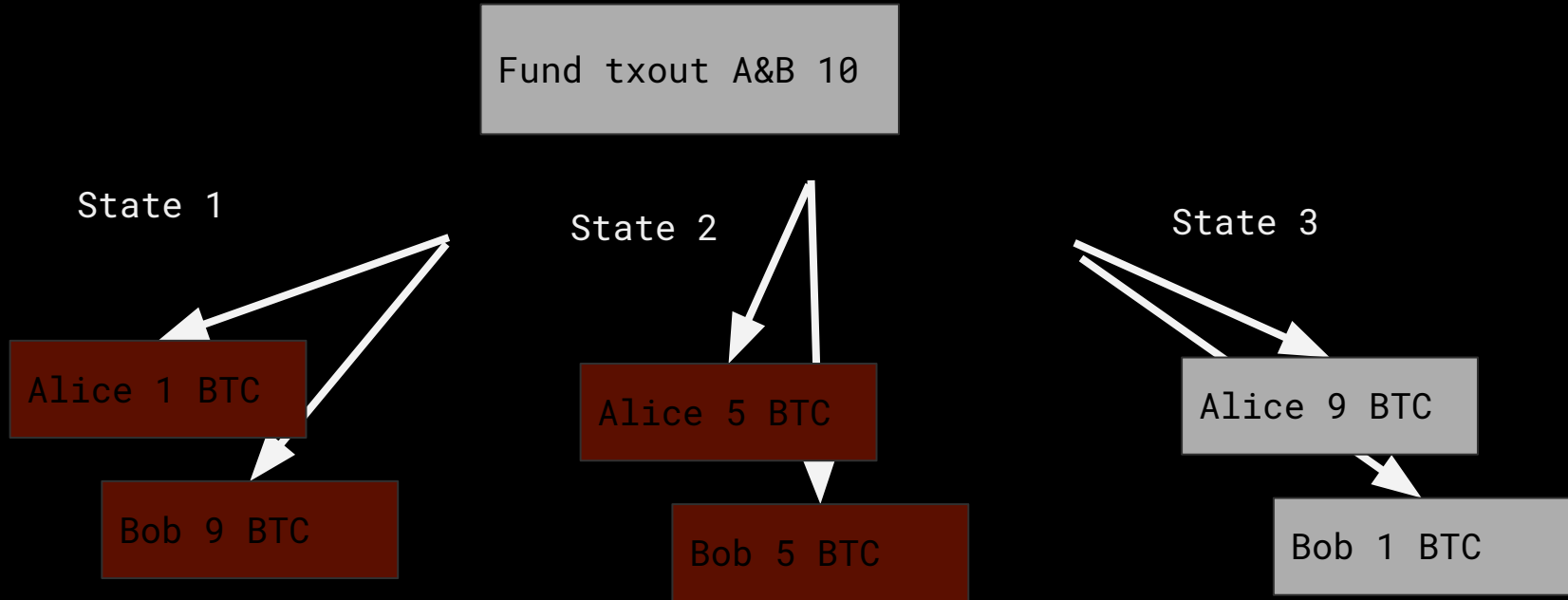
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signatures as private keys



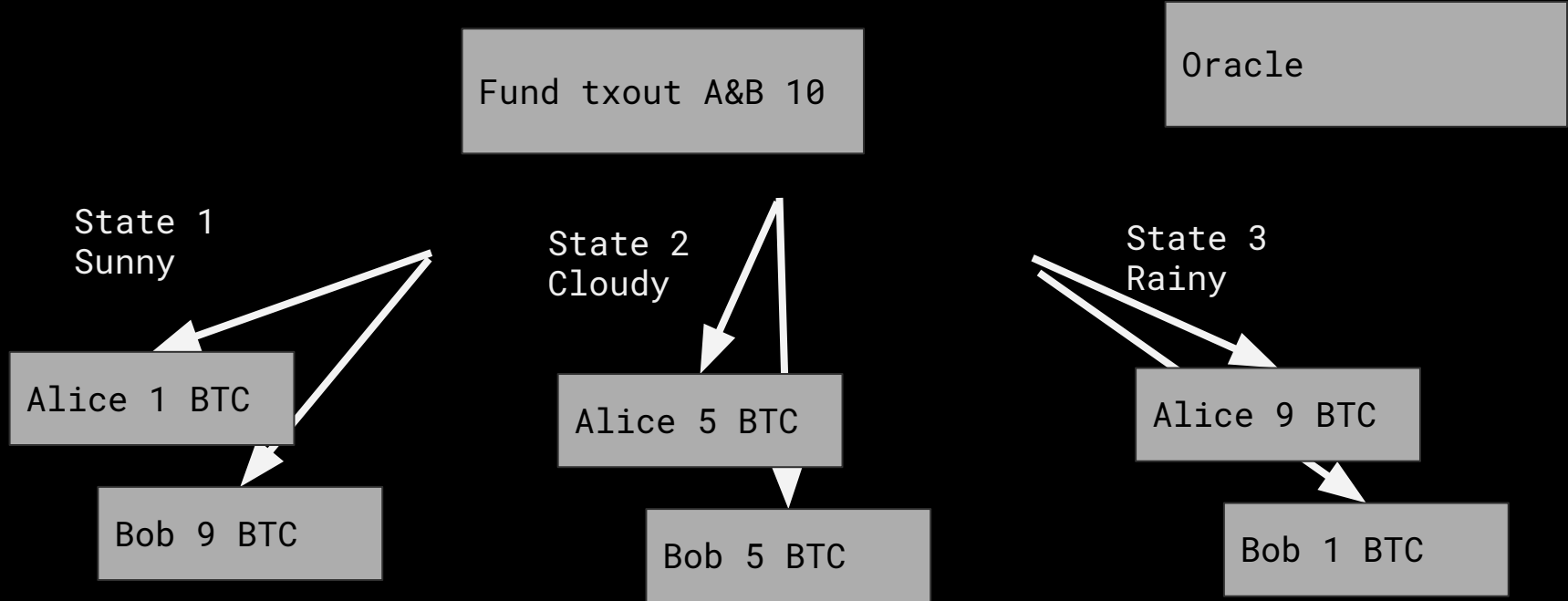
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signatures as private keys



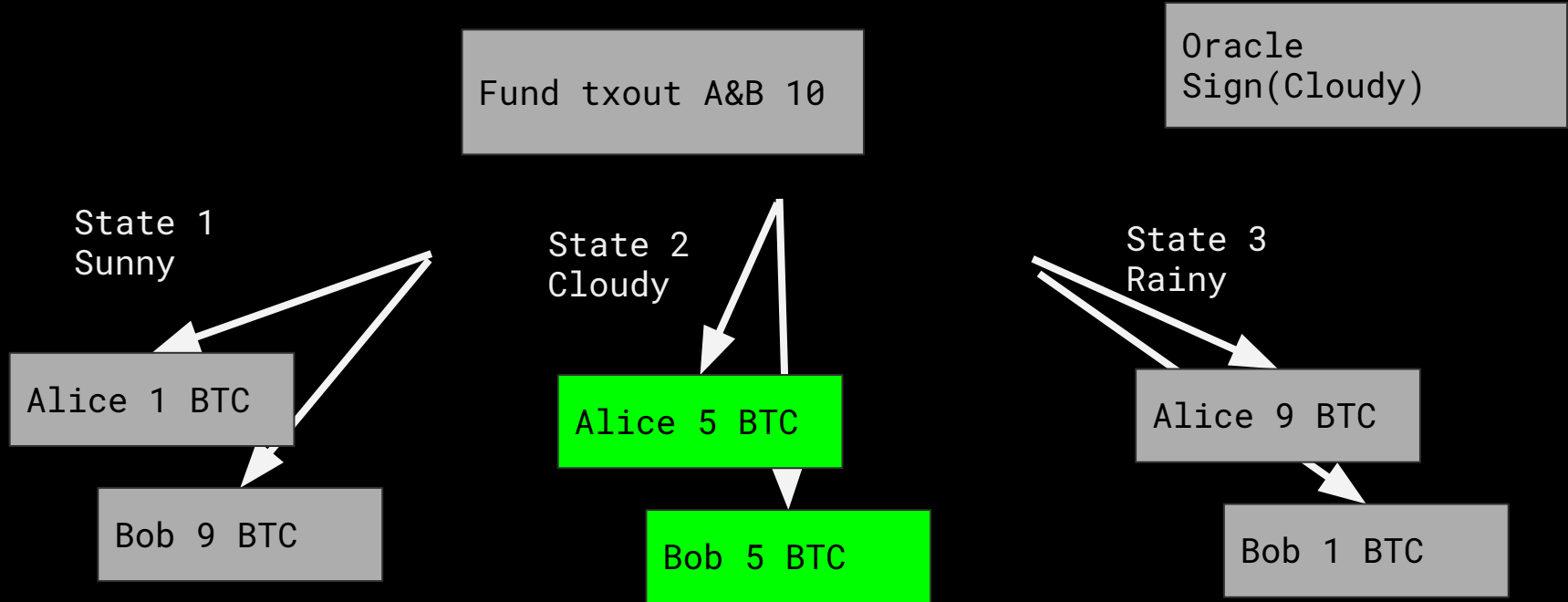
In Lightning, states are added sequentially, and validity is enforced by revealing private keys to previous states

signatures as private keys



In DLC all states are created at the start. Validity is determined by a non-interactive oracle signature.

signatures as private keys



In DLC all states are created at the start. Validity is determined by a non-interactive oracle signature.

Same script as LN

PubR OR (PubT AND time)

In lightning, The "correct" use is the
timeout, op_csv

In cases of fraud, the revocable key can
be used (half the key revealed)

In DLC, timeout is "incorrect", when
someone publishes the wrong tx.

time and DLCs

In LN, you need to always watch for fraud, as old states could be broadcast. Gotta grab that output.

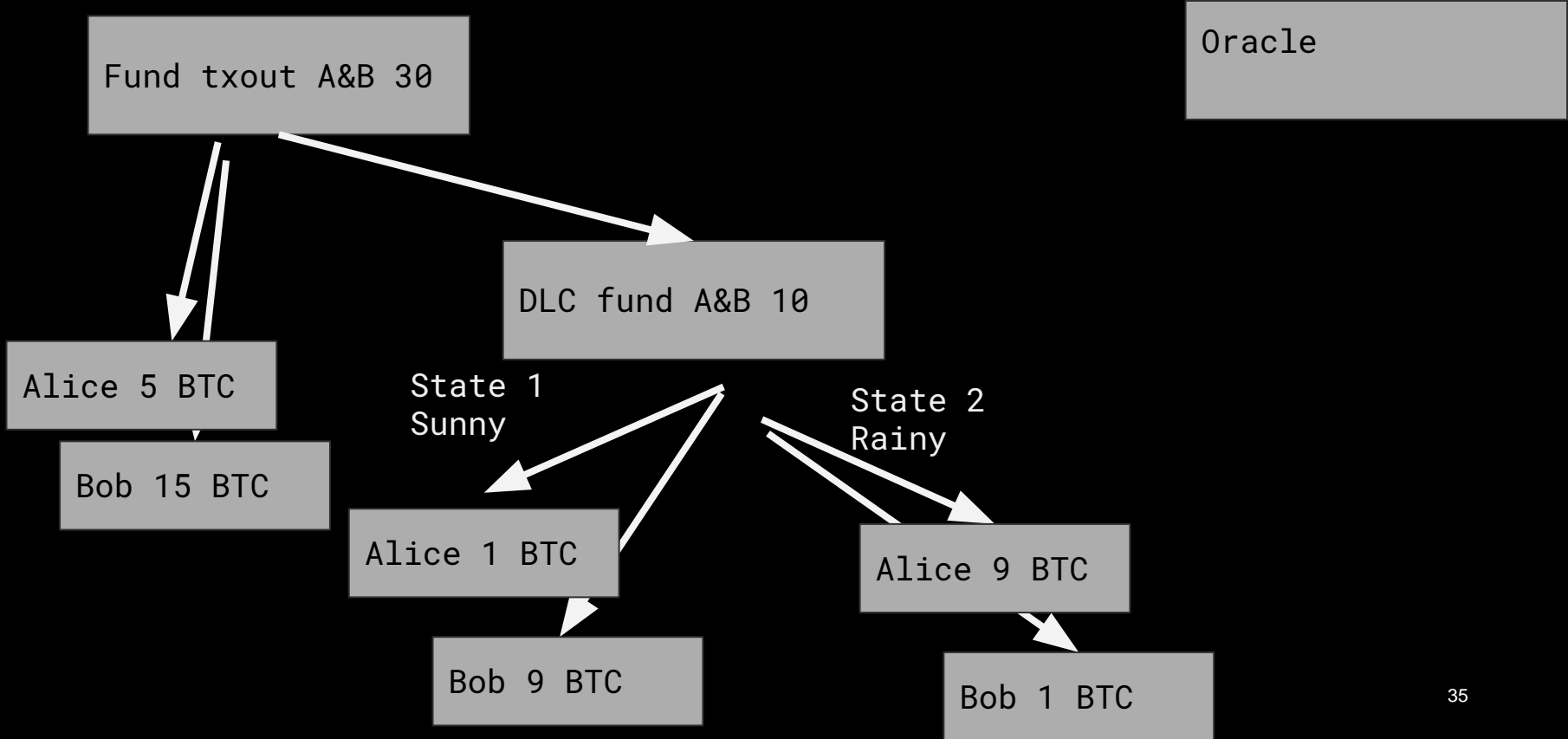
In DLC, you sweep the output as soon as you make it. Easier, and have the software broadcast both txs at the same time. No surprises.

DLCs within channels

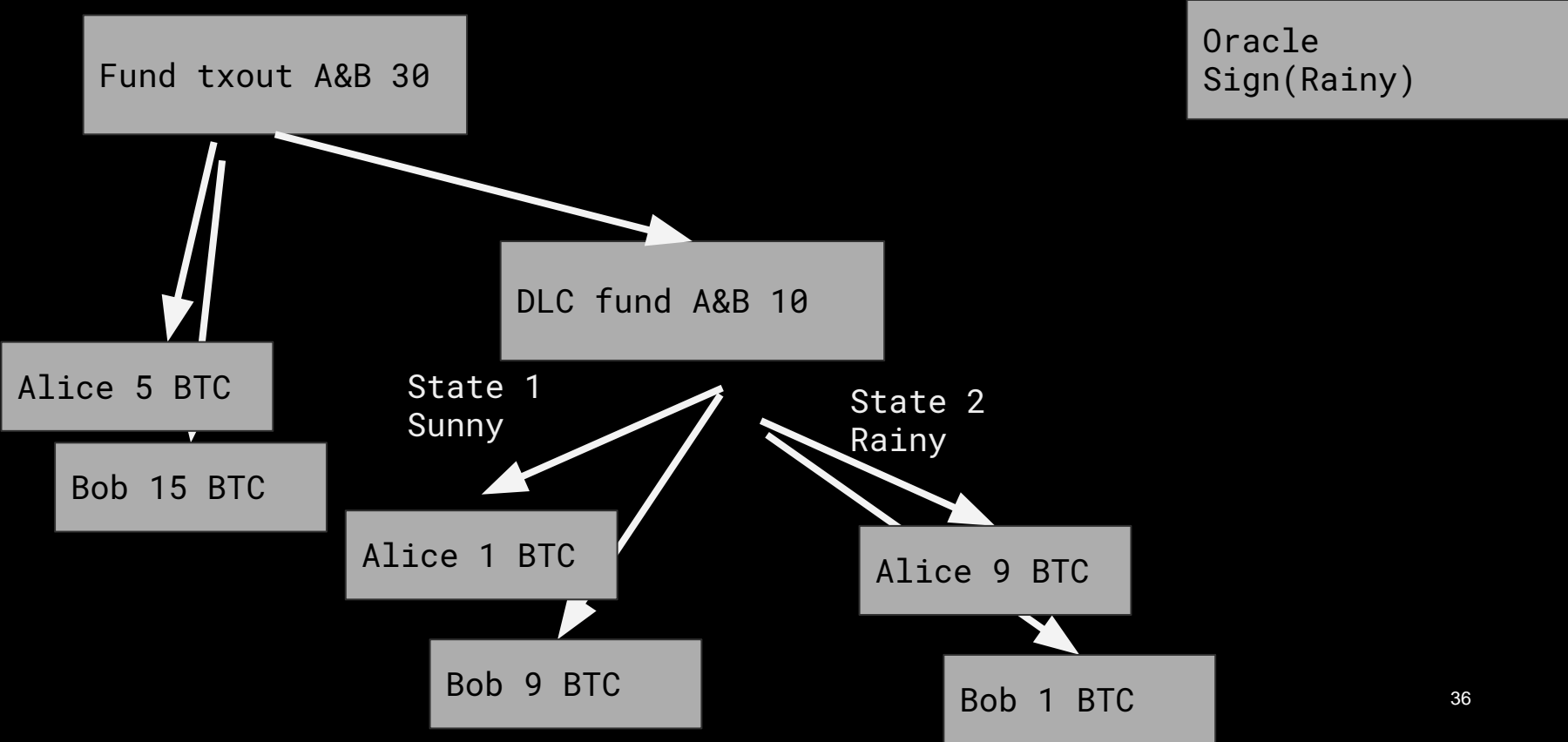
Make a DLC output from an LN channel

If parties cooperate, \emptyset txs get
broadcast to the blockchain

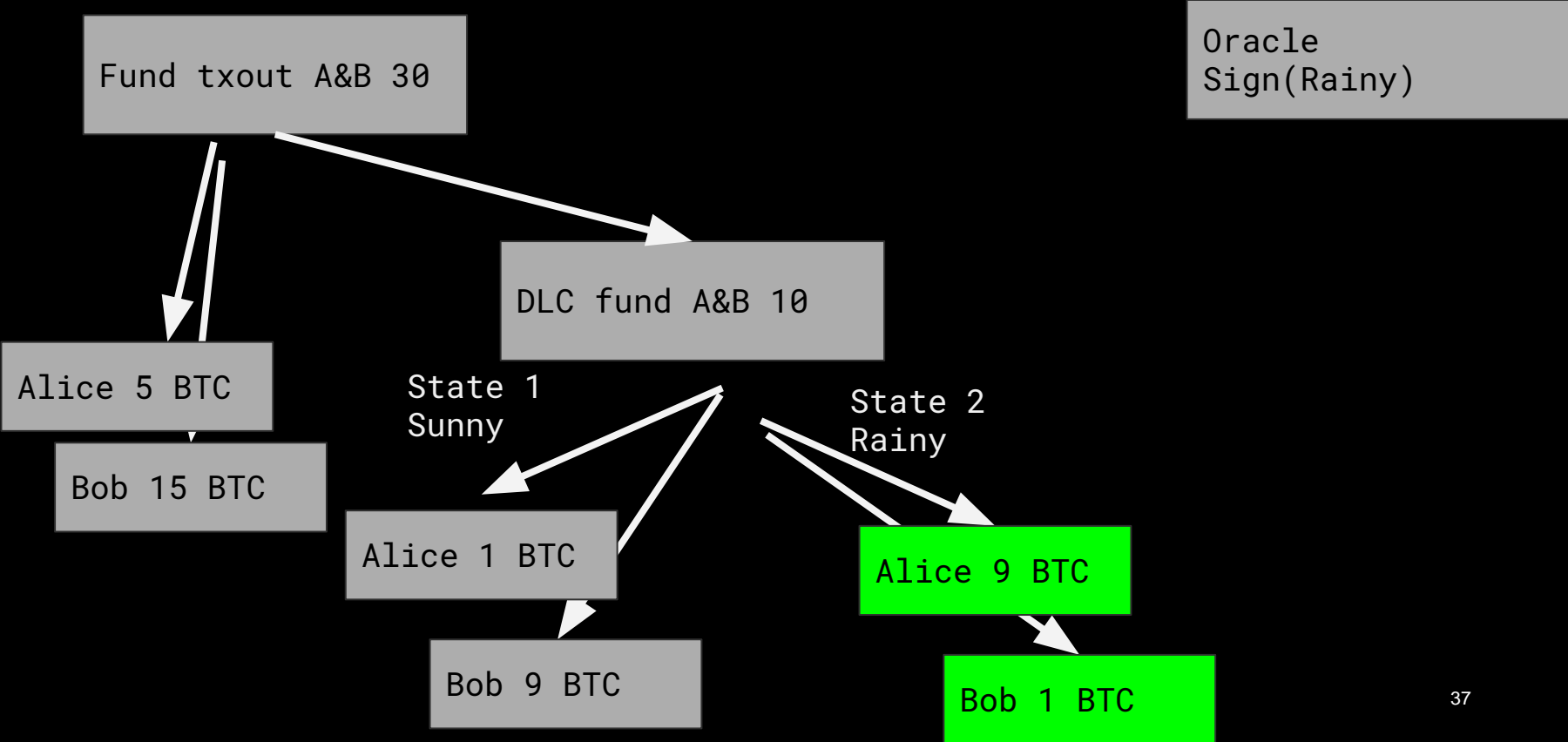
nested contracts



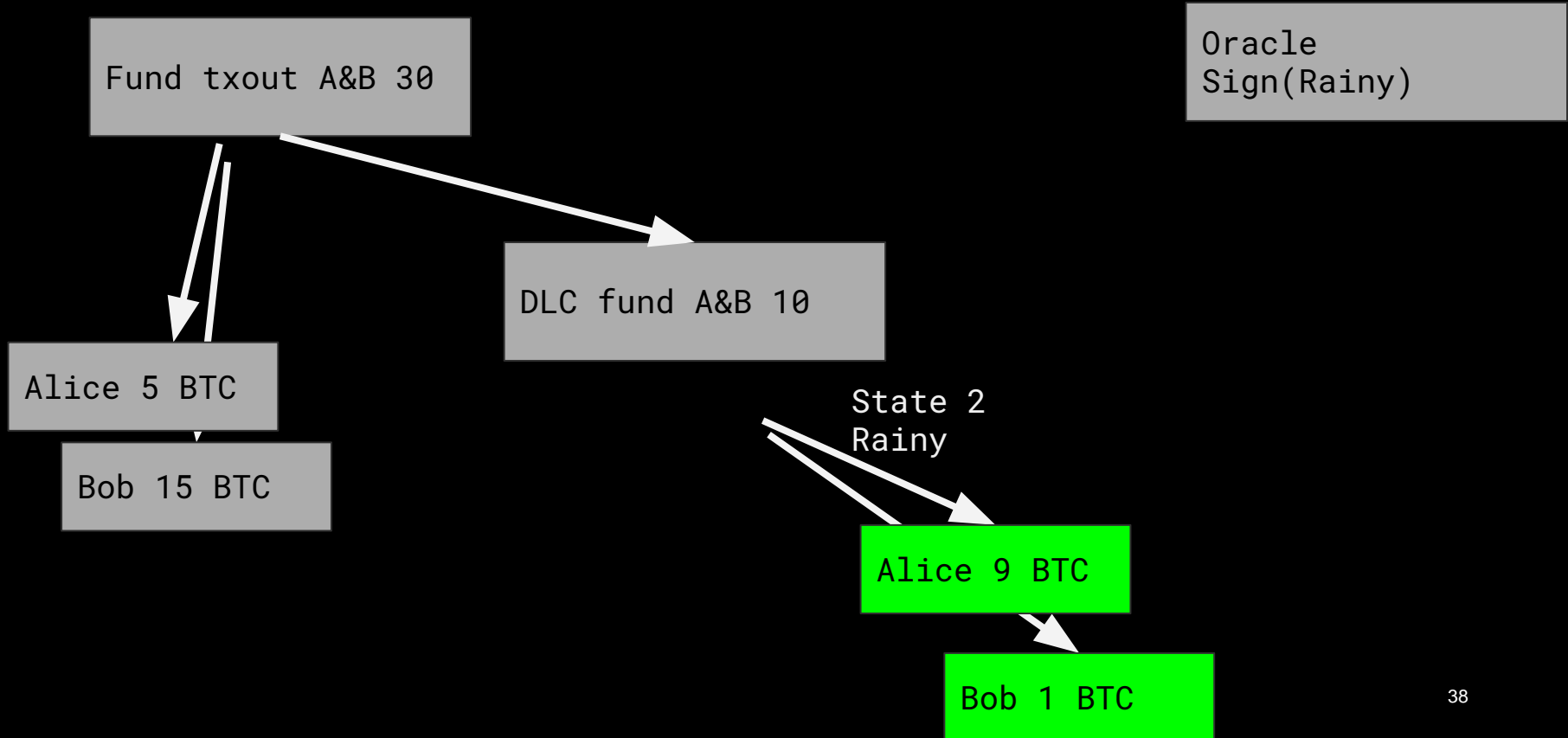
nested contracts



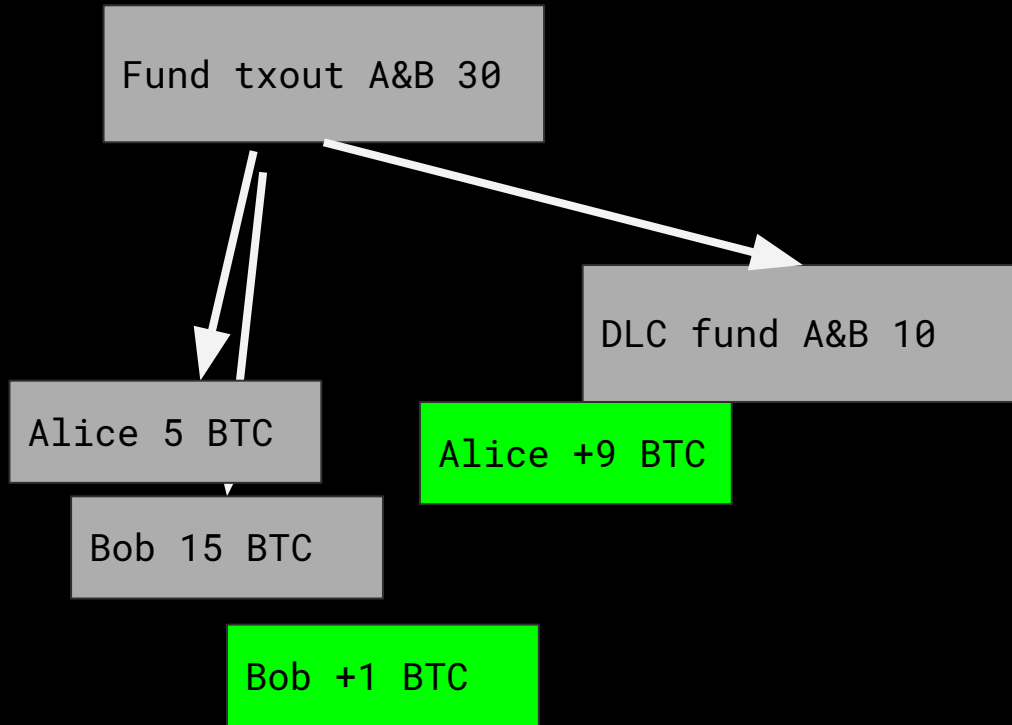
nested contracts



nested contracts

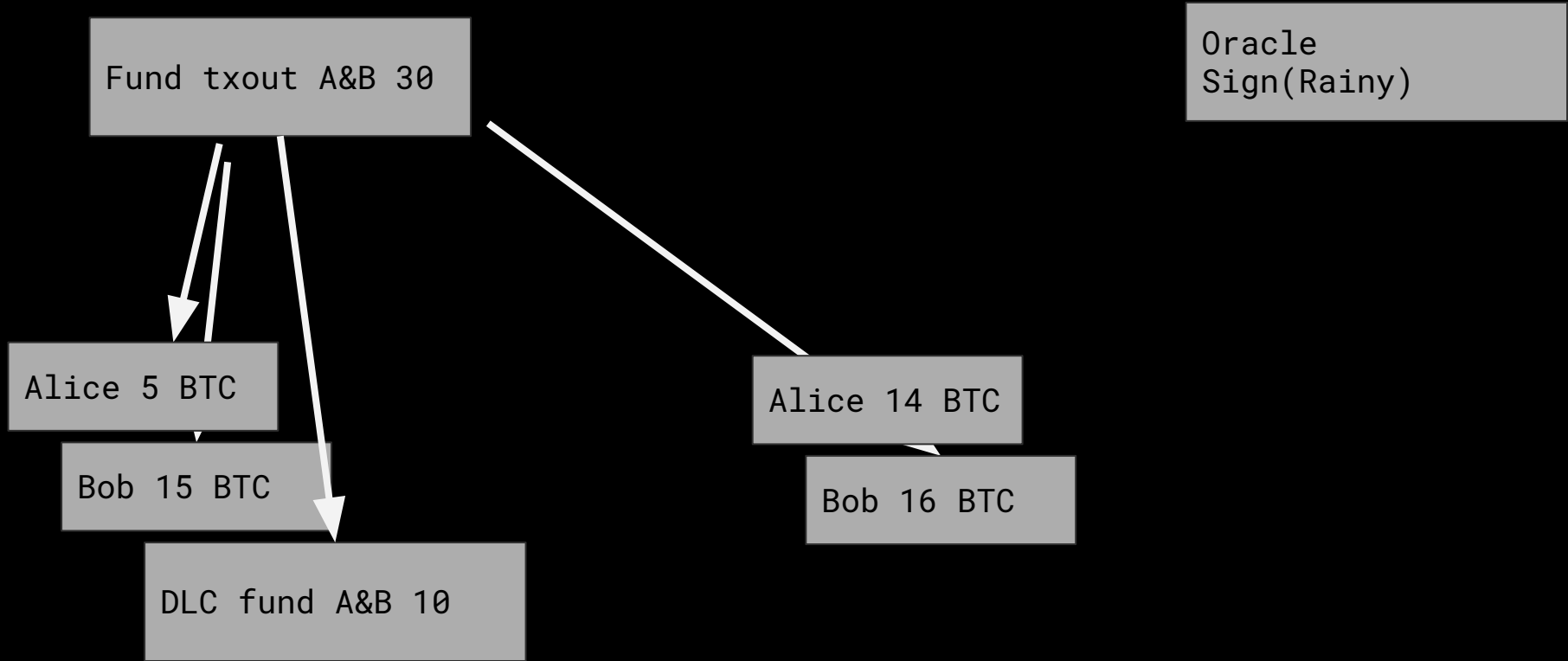


nested contracts

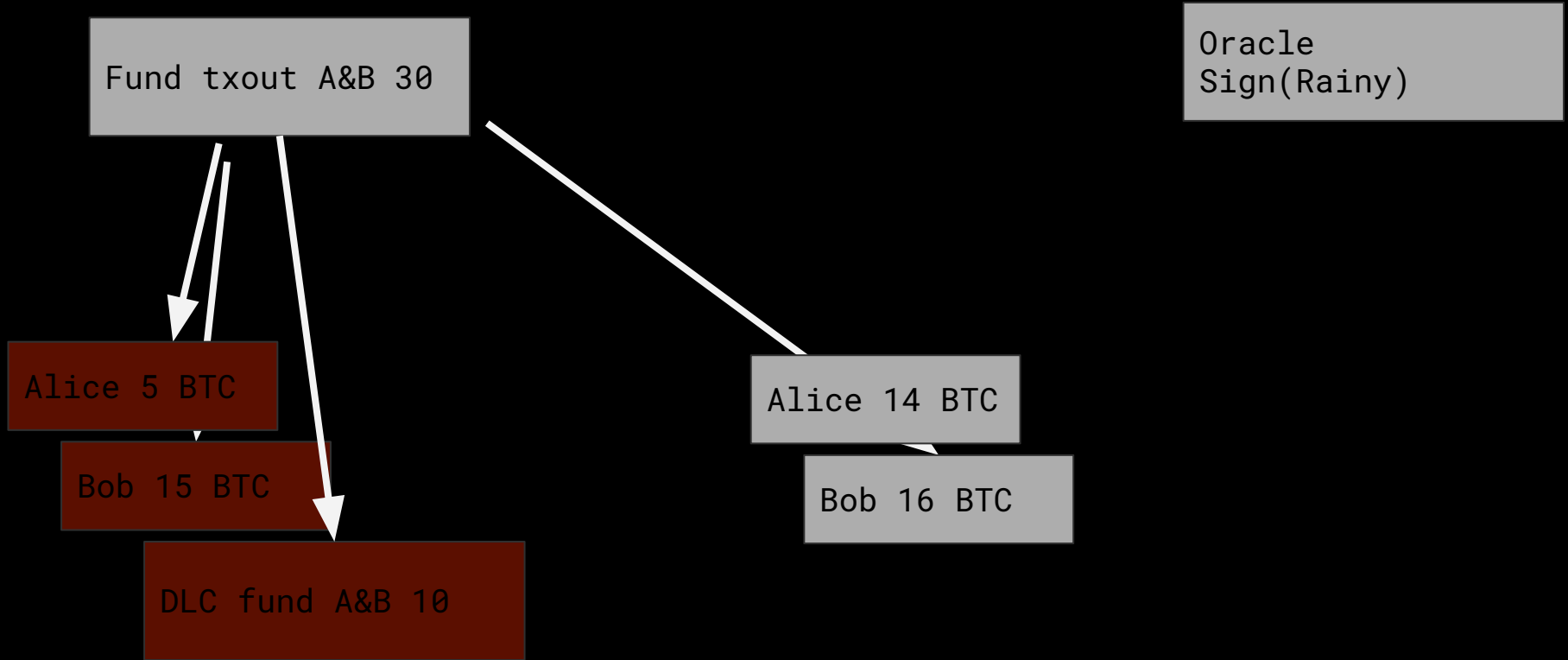


Oracle
Sign(Rainy)

nested contracts



nested contracts



DLC scalability

Can split the R value (and message) into a R-exponent and R-mantissa

Helps cut down the off-chain transactions needed in ranges which don't lead to different allocations

multi-oracle

Maybe Alice and Bob want to use 2 oracles. No problem.

$$s_a G + s_b G = s_c G$$

Just add the sG points. n of n , no size increase. (n of m , size blowup)

DLC use cases

Currency futures? Stocks?

Commodities? Sports? Insurance?

Pretty general; conditional payments based on any number or element from predetermined set.

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